



An Introduction to GPU Programming

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GPU Computing History

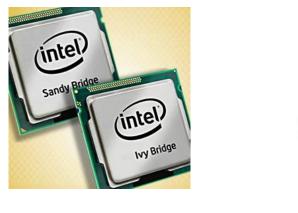
- The first GPU (Graphics Processing Unit)s were designed as graphics accelerators, supporting only specific fixed-function pipelines.
- Starting in the late 1990s, the hardware became increasingly programmable, culminating in NVIDIA's first GPU in 1999.
- Researchers were tapping its excellent floating point performance. The General Purpose GPU (GPGPU) movement had dawned.
- NVIDIA unveiled CUDA in 2006, the world's first solution for generalcomputing on GPUs.
- CUDA (Compute Unified Device Architecture) is a parallel computing platform and programming model created by NVIDIA and implemented by the GPUs that they produce.



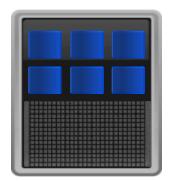


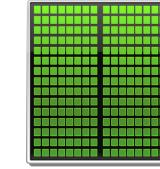


Add GPUs: Accelerate Science Applications CPU GPU













Why is GPU this different from a CPU?

Different goals produce different designs

- GPU assumes work load is highly parallel
- CPU must be good at everything, parallel or not
- > CPU: minimize latency experienced by 1 thread
 - big on-chip caches
 - sophisticated control logic

GPU: maximize throughput of all threads

- # threads in flight limited by resources => lots of resources (registers, bandwidth, etc.)
- multithreading can hide latency => skip the big caches
- share control logic across many threads







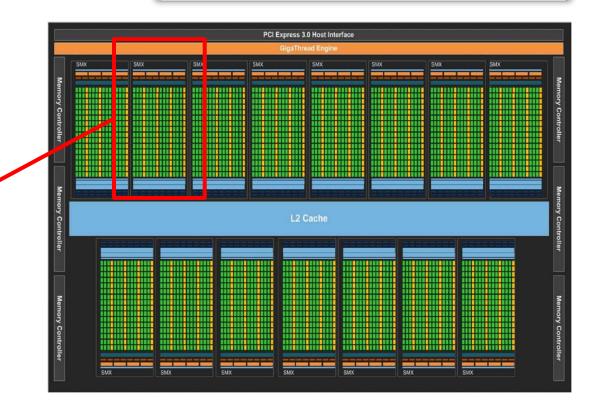
Overview of the GPU nodes

- CPU: Two 2.6 GHz 8-Core Sandy Bridge Xeon 64-bit Processors (16)
 - 64GB 1666MHz Ram
- GPU: Two NVIDIA Tesla K20Xm
 - 14 Streaming Multiprocessor (SMX)
 - 2688 SP Cores
 - 896 DP Cores
 - 6G global memory



SMX (192 SP, 64 DP)







Key Architectural Ideas



- SIMT (Single Instruction Multiple Thread) execution
 - threads run in groups of 32 called warps
 - threads in a warp share instruction unit (IU)
 - HW automatically handles divergence

Hardware multithreading

- HW resource allocation & thread scheduling
- HW relies on threads to hide latency

Threads have all resources needed to run

- any warp not waiting for something can run
- context switching is (basically) free

Instruction Cache			
Scheduler		Scheduler	
Dispatch		Dispatch	
Register File			
Core	Core	Core	Core
Core	Core	Core	Core
Core	Core	Core	Core
Core	Core	Core	Core
Core	Core	Core	Core
Core	Core	Core	Core
Core	Core	Core	Core
Core	Core	Core	Core
Load/Store Units x 16 Special Func Units x 4			
Interconnect Network			
64K Configurable Cache/Shared Mem			
Uniform Cache			







Enter CUDA

- Scalable parallel programming model
- Minimal extensions to familiar C/C++ environment
- Heterogeneous serial-parallel computing

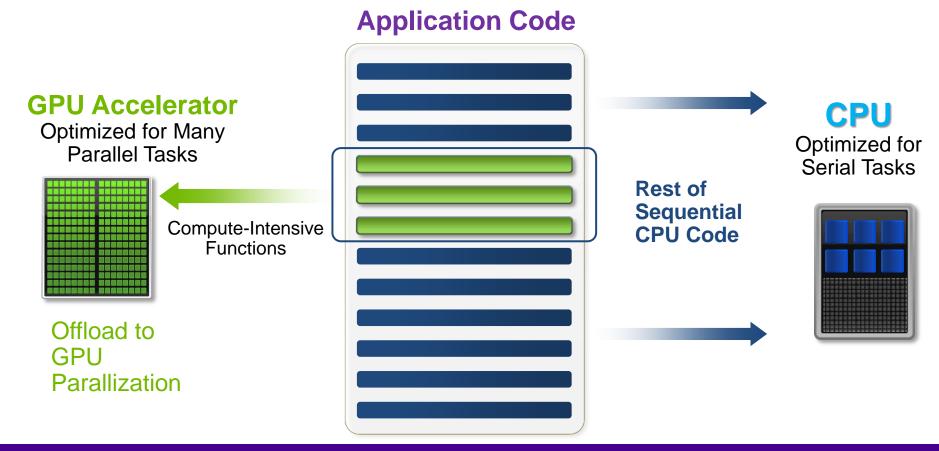






CUDA Execution Model

- Sequential code executes in a Host (CPU) thread
- Parallel code executes in many Device (GPU) threads across multiple processing elements



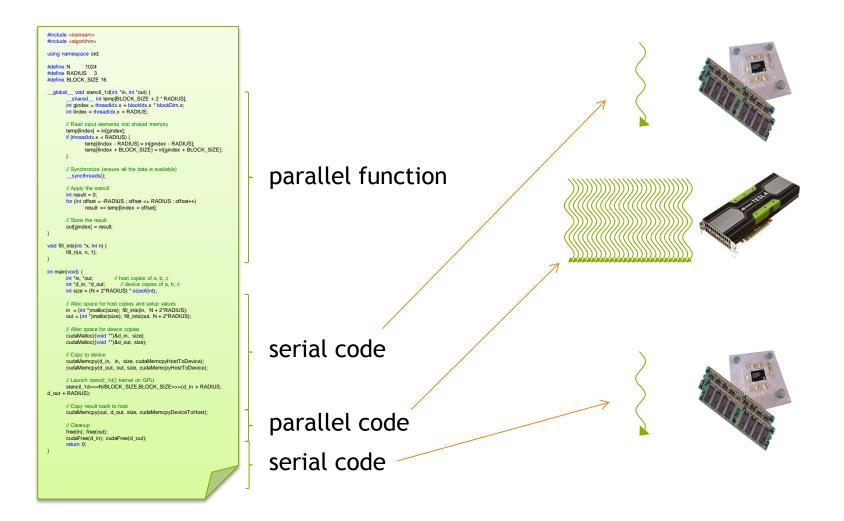
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Introduction to GPU Programming





Heterogeneous Computing





Introduction to GPU Programming

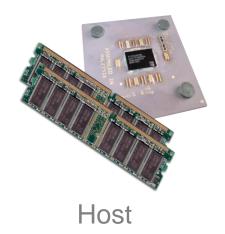




Heterogeneous Computing

Terminology:

- *Host* The CPU and its memory (host memory)
- Device The GPU and its memory (device memory)







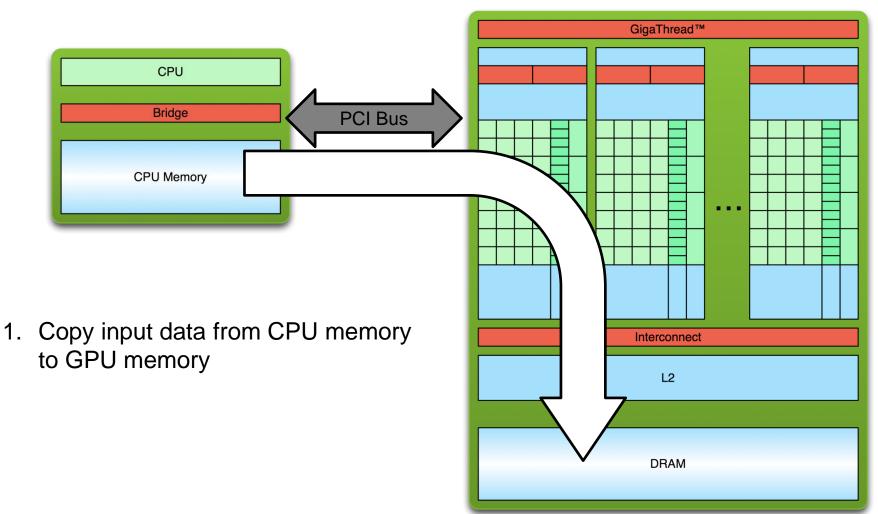


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Simple Processing Flow

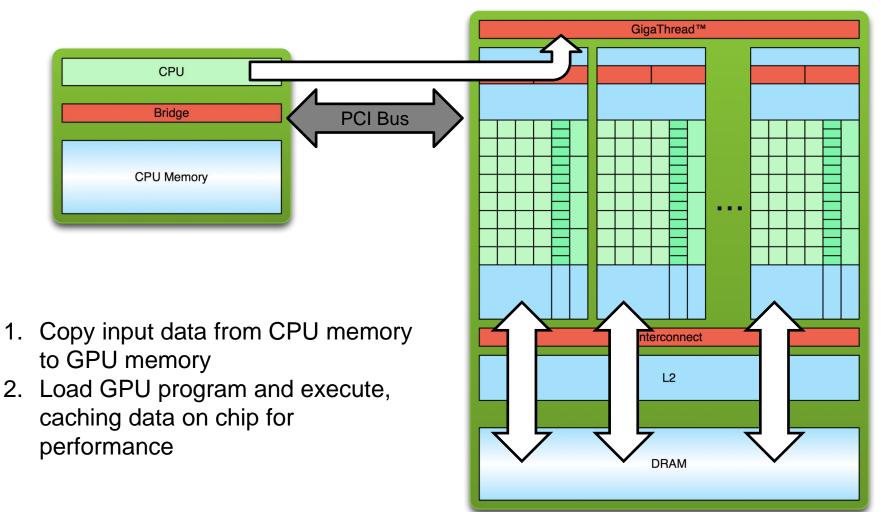








Simple Processing Flow

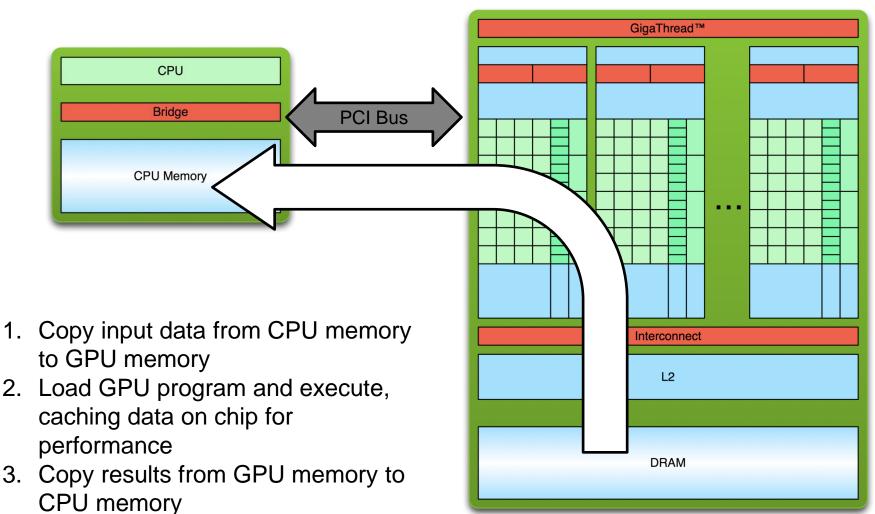








Simple Processing Flow

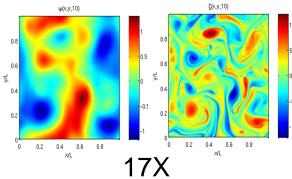


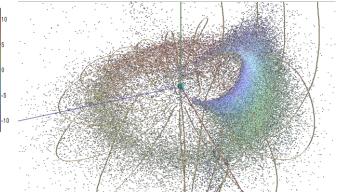




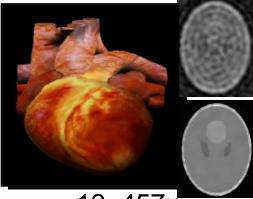


45X





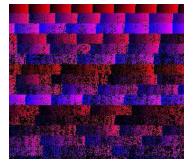




13–457x

110-240X

Motivation



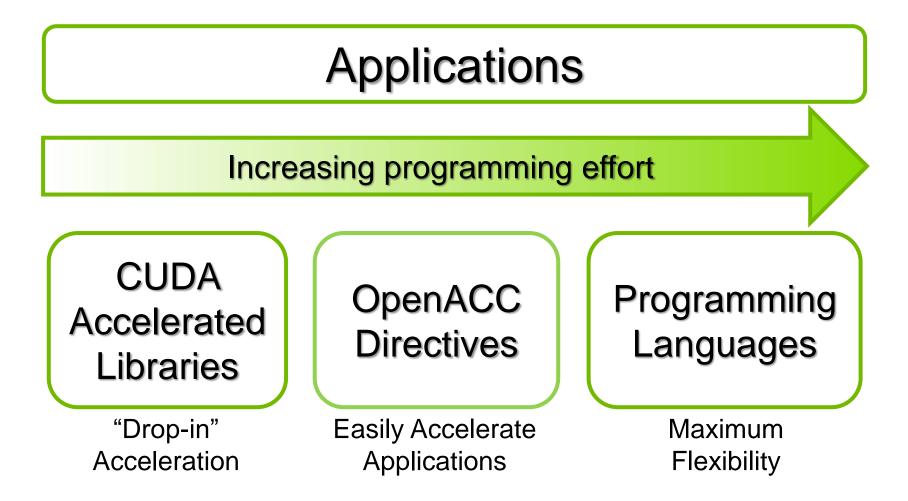


Introduction to GPU Programming





3 Ways to Accelerate Applications

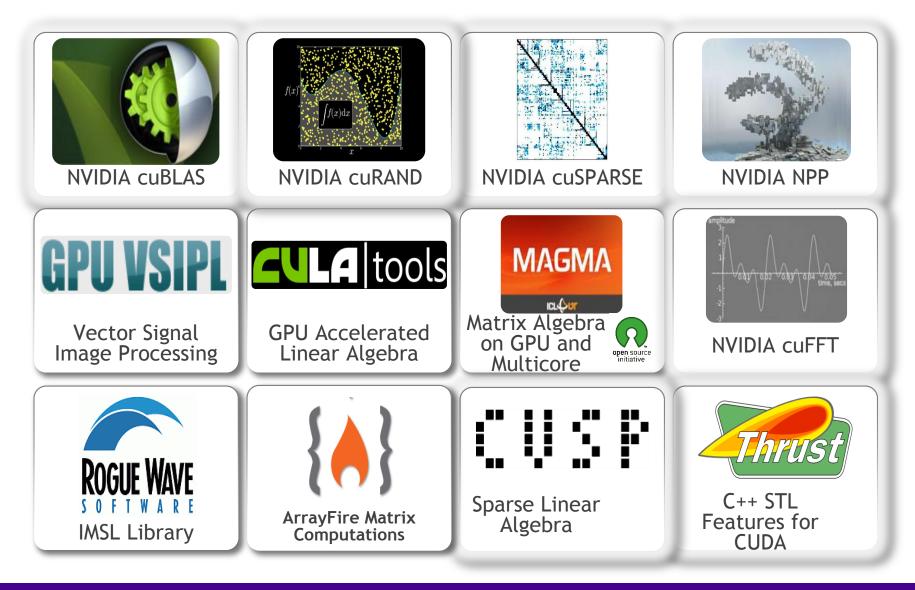








Some GPU-accelerated Libraries









GPU Programming Languages

Numerical analytics 🕨

MATLAB, Mathematica, LabVIEW

Fortran Þ

OpenACC, CUDA Fortran

OpenACC, CUDA C

C

Thrust, CUDA C++



F# 🕨

PyCUDA, Copperhead

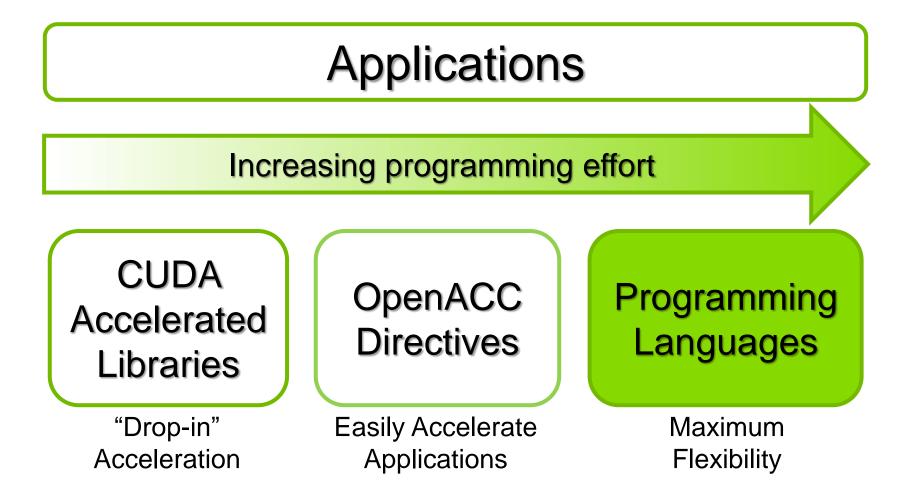
Alea.cuBase







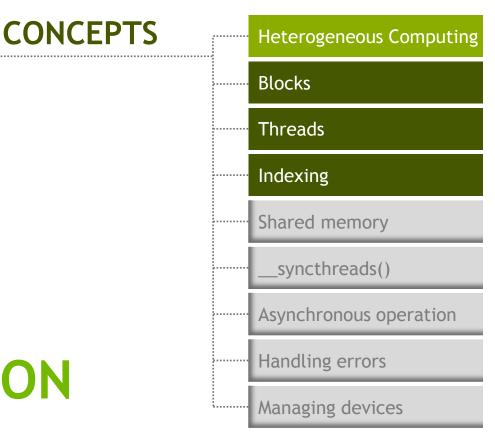
3 Ways to Accelerate Applications



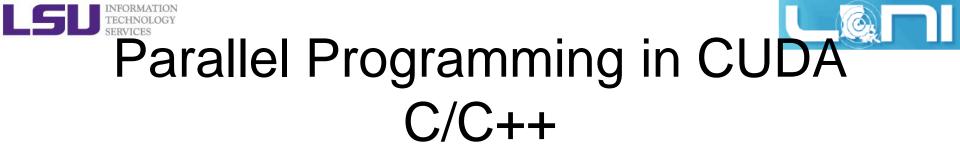




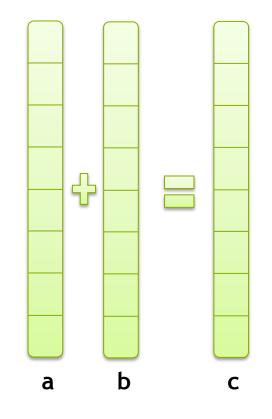




VECTOR ADDITION WITH CUDA



• We'll start by adding two integers and build up to vector addition









Addition on the Device

First recall how to write a pure C function:

```
void add(int *a, int *b, int *c) {
     *c = *a + *b;
}
```

> Then we have a simple kernel to add two integers

```
__global___void add(int *a, int *b, int *c) {
    *c = *a + *b;
}
```

As before global is a CUDA C/C++ keyword meaning
 add() is a kernel function that will execute on the device
 add() will be called from the host







Addition on the Device

> Note that we use pointers for the variables

```
__global__ void add(int *a, int *b, int *c) {
    *c = *a + *b;
}
```

- add() runs on the device, so a, b and c must point to device memory
- > We need to allocate memory on the GPU







Memory Management

> Host and device memory are separate entities

- Device pointers point to GPU memory May be passed to/from host code
 May not be dereferenced in host code
- Host pointers point to CPU memory
 May be passed to/from device code
 May not be dereferenced in device code





Simple CUDA API for handling device memory

- cudaMalloc(),cudaFree(),cudaMemcpy()
- Similar to the C equivalents malloc(), free(), memcpy()





Addition on the Device: add()

Returning to our add () kernel

```
__global__ void add(int *a, int *b, int *c) {
    *c = *a + *b;
}
```

Let's take a look at main()...







Addition on the Device: main()

int main(void) {
 int a, b, c; // host copies of a, b, c
 int *d_a, *d_b, *d_c; // device copies of a, b, c
 int size = sizeof(int);

// Allocate space for device copies of a, b, c
cudaMalloc((void **)&d_a, size);
cudaMalloc((void **)&d_b, size);
cudaMalloc((void **)&d_c, size);

// Setup input values
a = 2;
b = 7;







Addition on the Device: main()

// Copy inputs to device
cudaMemcpy(d_a, &a, size, cudaMemcpyHostToDevice);
cudaMemcpy(d b, &b, size, cudaMemcpyHostToDevice);

// Launch add() kernel on GPU
add<<<1,1>>>(d_a, d_b, d_c);

// Copy result back to host
cudaMemcpy(&c, d_c, size, cudaMemcpyDeviceToHost);

```
// Cleanup
```

```
cudaFree(d_a); cudaFree(d_b); cudaFree(d_c);
return 0;
```

}





Compile and Run

Changes to the ~/.soft file:

```
[fchen14@mike2 gpuex]$ cat ~/.soft
+cuda-5.5.22
+Intel-13.1.3
+portland-14.3
@default
```

Request an interactive session in GPU queue:

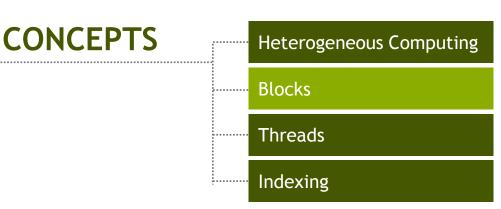
```
qsub -I -X -l nodes=1:ppn=16 -l walltime=01:00:00 -q gpu -A
your_allocation_name
```

Compile and run the first vector addition:

```
[fchen14@mike424 gpuex]$ nvcc my_vec_add.cu
[fchen14@mike424 gpuex]$ ./a.out
c=9
```







RUNNING IN PARALLEL



Introduction to GPU Programming





Moving to Parallel

- GPU computing is about massive parallelism
 - So how do we run code in parallel on the device?

```
add<<< 1, 1 >>>();
```

```
add<<< N, 1 >>>();
```

> Instead of executing add() once, execute N times in parallel







Vector Addition on the Device

- With add() running in parallel we can do vector addition
- > Terminology: each parallel invocation of add() is referred to as a block
 - The set of blocks is referred to as a grid
 - Each invocation can refer to its block index using **blockIdx**.x

```
__global___void add(int *a, int *b, int *c) {
    c[blockIdx.x] = a[blockIdx.x] + b[blockIdx.x];
}
```

By using blockIdx.x to index into the array, each block handles a different index





Vector Addition on the Device

```
__global___void add(int *a, int *b, int *c) {
    c[blockIdx.x] = a[blockIdx.x] + b[blockIdx.x];
}
```

> On the device, each block can execute in parallel:









Vector Addition on the Device: add()

> Returning to our parallelized add() kernel

```
__global___void add(int *a, int *b, int *c) {
    c[blockIdx.x] = a[blockIdx.x] + b[blockIdx.x];
}
```

Let's take a look at main()...







Vector Addition on the Device: main()

#define N 512

```
int main(void) {
```

```
int *a *b *c // host copies of a, b, c
int *d_a, *d_b, *d_c; // device copies of a, b, c
int size = N * sizeof(int);
```

```
// Alloc space for device copies of a, b, c
cudaMalloc((void **)&d_a, size);
cudaMalloc((void **)&d_b, size);
cudaMalloc((void **)&d_c, size);
```

```
// Alloc space for host copies of a, b, c and setup input values
a = (int *)malloc(size);
b = (int *)malloc(size);
c = (int *)malloc(size);
```







Vector Addition on the Device: main()

// Copy inputs to device
cudaMemcpy(d_a, a, size, cudaMemcpyHostToDevice);
cudaMemcpy(d_b, b, size, cudaMemcpyHostToDevice);

// Launch add() kernel on GPU with N blocks
add<<<N,1>>>(d_a, d_b, d_c);

// Copy result back to host
cudaMemcpy(c, d_c, size, cudaMemcpyDeviceToHost);

```
// Cleanup
free(a); free(b); free(c);
cudaFree(d_a); cudaFree(d_b); cudaFree(d_c);
return 0;
```



}





Review (1 of 2)

- Difference between host and device
 - Host CPU
 - *Device* GPU
- Using __global___ to declare a function as device code
 - Executes on the device
 - Called from the host
- Passing parameters from host code to a device function







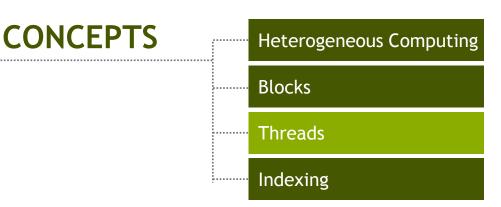
Review (2 of 2)

- Basic device memory management
 - cudaMalloc()
 - cudaMemcpy()
 - cudaFree()
- Launching parallel kernels
 - Launch \mathbf{N} copies of add() with add <<< \mathbf{N} , 1>>> (...);
 - Use **blockIdx**.**x** to access block index









INTRODUCING THREADS



Introduction to GPU Programming





CUDA Threads

- Ferminology: a block can be split into parallel threads
 - OR: block is composed of threads
- > Let's change add() to use parallel threads instead of parallel blocks
 ___global___ void add(int *a, int *b, int *c) {
 c[threadIdx.x] = a[threadIdx.x] + b[threadIdx.x];
 }
- We use threadIdx.x instead of blockIdx.x
- Need to make one change in main()...







Vector Addition Using Threads: main()

```
#define N 512
int main(void) {
   int *a, *b, *c;
                         // host copies of a, b, c
   int *d_a, *d_b, *d_c; // device copies of a, b, c
   int size = N * sizeof(int);
   // Alloc space for device copies of a, b, c
   cudaMalloc((void **)&d a, size);
   cudaMalloc((void **)&d b, size);
   cudaMalloc((void **)&d c, size);
   // Alloc space for host copies of a, b, c and setup input values
   a = (int *)malloc(size);
   b = (int *)malloc(size);
   c = (int *)malloc(size);
```

for (int i=0; i<N; i++) a[i]=2, b[i]=7;</pre>





Vector Addition Using Threads: main()

// Copy inputs to device
cudaMemcpy(d_a, a, size, cudaMemcpyHostToDevice);
cudaMemcpy(d_b, b, size, cudaMemcpyHostToDevice);

// Launch add() kernel on GPU with N threads
add<<<1,N>>>(d_a, d_b, d_c);

// Copy result back to host
cudaMemcpy(c, d_c, size, cudaMemcpyDeviceToHost);

// Cleanup

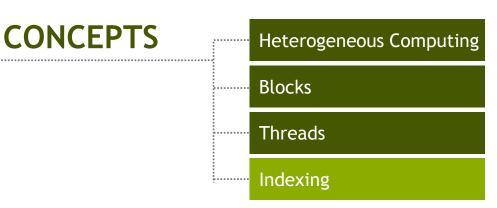
```
free(a); free(b); free(c);
cudaFree(d_a); cudaFree(d_b); cudaFree(d_c);
return 0;
```



}







COMBINING THREADS AND BLOCKS





Combining Blocks and Threads

- We've seen parallel vector addition using:
 - Many blocks with one thread each
 - One block with many threads
- Let's adapt vector addition to use both blocks and threads
- > Why? We'll come to that...
- First let's discuss data indexing...







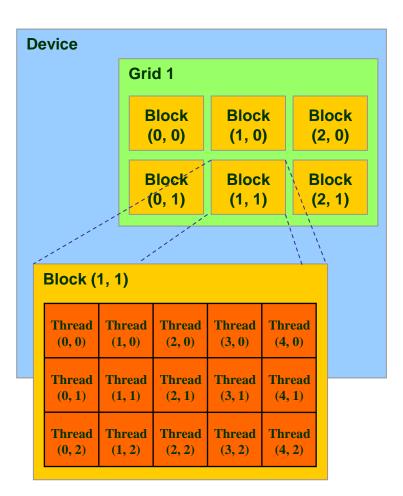
IDs and Dimensions

> Threads:

- 3D IDs, unique within a block

> Blocks:

- 2D IDs, unique within a grid
- Dimensions set at launch
 - Can be unique for each grid
- Built-in variables:
 - threadIdx, blockIdx
 - blockDim, gridDim
- We will only discuss the usage of one dimension (x)



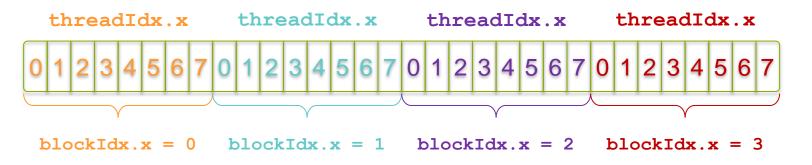






Indexing Arrays with Blocks and Threads

- > No longer as simple as using blockldx.x and threadldx.x
 - Consider indexing an array with one element per thread (8 threads/block)



With M (M=8 here) threads per block a unique index for each thread is given by:

```
int index = threadIdx.x + blockIdx.x * M;
```

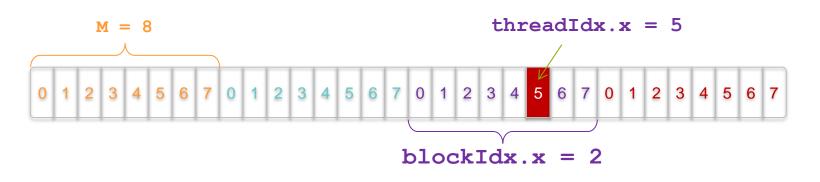




Indexing Arrays: Example

Which thread will operate on the red element?

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31



int index = threadIdx.x + blockIdx.x * M; = 5 + 2 * 8; = 21;







Vector Addition with Blocks and Threads

- > Use the built-in variable blockDim.x for threads per block int index = threadIdx.x + blockIdx.x * blockDim.x;
- > Combined version of add() to use parallel threads and parallel blocks

```
__global___void add(int *a, int *b, int *c) {
    int index = threadIdx.x + blockIdx.x * blockDim.x;
    c[index] = a[index] + b[index];
}
```

> What changes need to be made in main()?







Addition with Blocks and Threads:

main()

```
#define N (2048*2048)
#define THREADS_PER_BLOCK 512
int main(void) {
    int *a, *b, *c; // host copies of a, b, c
    int *d_a, *d_b, *d_c; // device copies of a, b, c
    int size = N * sizeof(int);
```

```
// Alloc space for device copies of a, b, c
cudaMalloc((void **)&d_a, size);
cudaMalloc((void **)&d_b, size);
cudaMalloc((void **)&d c, size);
```

```
// Alloc space for host copies of a, b, c and setup input values
a = (int *)malloc(size);
b = (int *)malloc(size);
c = (int *)malloc(size);
```







Addition with Blocks and Threads: main()

// Copy inputs to device

cudaMemcpy(d_a, a, size, cudaMemcpyHostToDevice); cudaMemcpy(d_b, b, size, cudaMemcpyHostToDevice);

// Launch add() kernel on GPU
add<<<N/THREADS_PER_BLOCK, THREADS_PER_BLOCK>>>(d_a, d_b, d_c);

// Copy result back to host
cudaMemcpy(c, d_c, size, cudaMemcpyDeviceToHost);

// Cleanup

```
free(a); free(b); free(c);
cudaFree(d_a); cudaFree(d_b); cudaFree(d_c);
return 0;
```

}





Handling Arbitrary Vector Sizes

- Typical problems are not friendly multiples of blockDim.x
- Avoid accessing beyond the end of the arrays:

```
__global___void add(int *a, int *b, int *c, int n) {
    int index = threadIdx.x + blockIdx.x * blockDim.x;
    if (index < n)
        c[index] = a[index] + b[index];
}</pre>
```

Update the kernel launch:

add<<<(N + M-1) / M,M >>>(d_a, d_b, d_c, N);





Review



Launching parallel kernels

- Launch \mathbf{N} copies of add() with add<<< $\mathbf{N}/\mathbf{M}, \mathbf{M}$ >>>(...);
- Use **blockIdx**. **x** to access block index
- Use **threadIdx**.**x** to access thread index within block
- Allocate elements to threads:

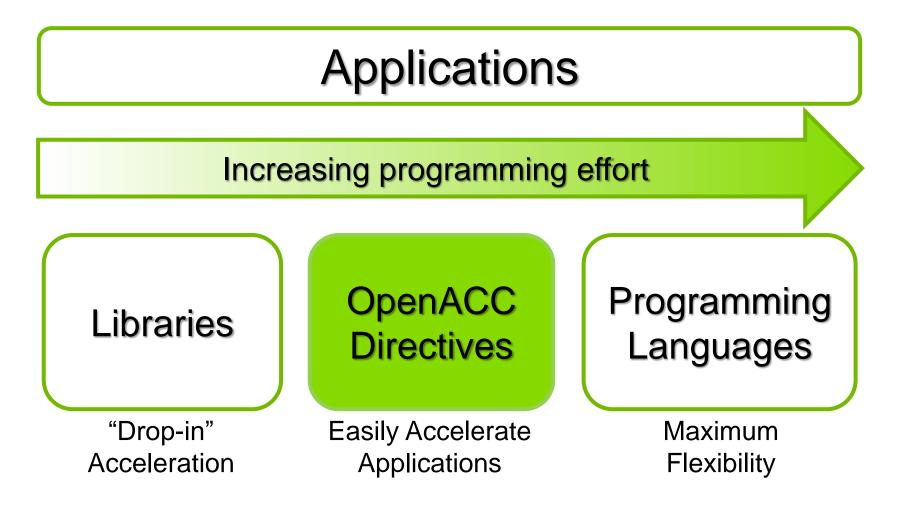
int index = threadIdx.x + blockIdx.x * blockDim.x;







3 Ways to Accelerate Applications









To be covered

- > OpenACC overview
- First OpenACC program and basic OpenACC directives
- Data region concept
- How to parallize our examples:
 - Laplacian solver
- Hands-on exercise
 - Matrix Multiplication
 - SAXPY
 - Calculate π







What is OpenACC

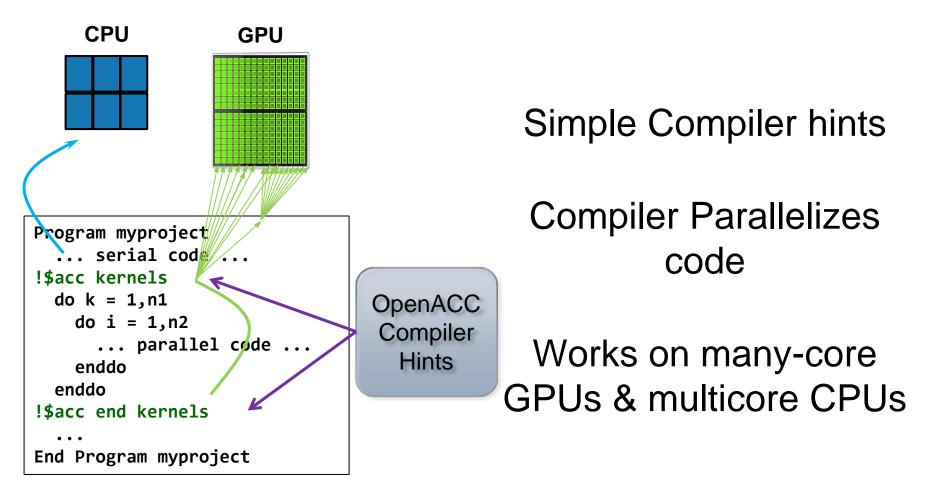
- OpenACC (for Open Accelerators) is a programming standard for parallel computing developed by Cray, CAPS, Nvidia and PGI. The standard is designed to simplify parallel programming of heterogeneous CPU/GPU systems.
- It provides a model for accelerator programming that is portable across operating systems and various types of host CPUs and accelerators.
- Full OpenACC 2.0 Specification available online
 - http://www.openacc-standard.org/
 - Implementations available now from PGI, Cray, and CAPS







OpenACC Directives



Your original Fortran or C code





The Standard for GPU Directives

Simple and high-level :

- Directive are the easy path to accelerate compute intensive applications. Non-GPU programmers can play along.
- Single Source: Compile the same program for accelerators or serial, No involvement of OpenCL, CUDA, etc.

Open and performance portable:

- OpenACC is an open GPU directives standard, making GPU programming straightforward and portable across parallel and multi-core processors
- Supports GPU accelerators and co-processors from multiple vendors, current and future versions.

Powerful and Efficient:

- Directives allow complete access to the massive parallel power of GPU.
- Experience shows very favorable comparison to low-level implementations of same algorithms.
- Developers can port and tune parts of their application as resources and profiling dictates. No need to restructure the program.





Directive-based programming

Directives provide a high-level alternative

- Based on original source code (Fortran, C, C++)
- Easier to maintain/port/extend code
- Users with OpenMP experience find it a familiar programming model
- Compiler handles repetitive coding (cudaMalloc, cudaMemcpy...)
- Compiler handles default scheduling; user tunes only where needed

Possible performance sacrifice

- Small performance sacrifice is acceptable
- trading-off portability and productivity against this
- after all, who hand-codes in assembly for CPUs these days?

As researchers in science and engineering, you often need to balance between:

- □ Time needed to develop your code
- **Time needed to focus on the problem itself**





General Directive Syntax and Scope

Fortran

```
!$acc directive [clause [,] clause]...]
```

```
Often paired with a matching end directive surrounding a structured code block
```

!\$acc end directive

```
≻ C
```

```
#pragma acc directive [clause [,] clause]...]
{
   Often followed by a structured code block (compound
   statement)
}
```







The "restrict" keyword in C

- Declaration of intent given by the programmer to the compiler
 - Applied to a pointer, e.g. float *restrict ptr;
 - Meaning: "for the lifetime of ptr, only it or a value directly derived from it (such as ptr + 1) will be used to access the object to which it points"*
 - In simple, the ptr will only point to the memory space of itself
- > OpenACC compilers often require restrict to determine independence.
 - Otherwise the compiler can't parallelize loops that access ptr
 - Note: if programmer violates the declaration, behavior is undefined.

THE RESTRICT CONTRACT

I, [insert your name], a PROFESSIONAL or AMATEUR [circle one] programmer, solemnly declare that writes through this pointer will not effect the values read through any other pointer available in the same context which is also declared as restricted.

* Your agreement to this contract is implied by use of the restrict keyword ;)

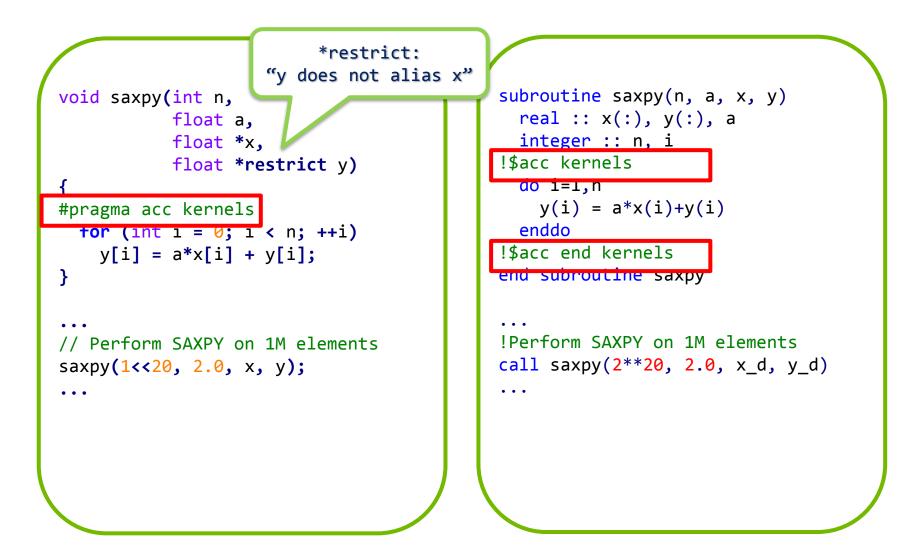
*http://en.wikipedia.org/wiki/Restrict







The First Simple Exercise: SAXPY



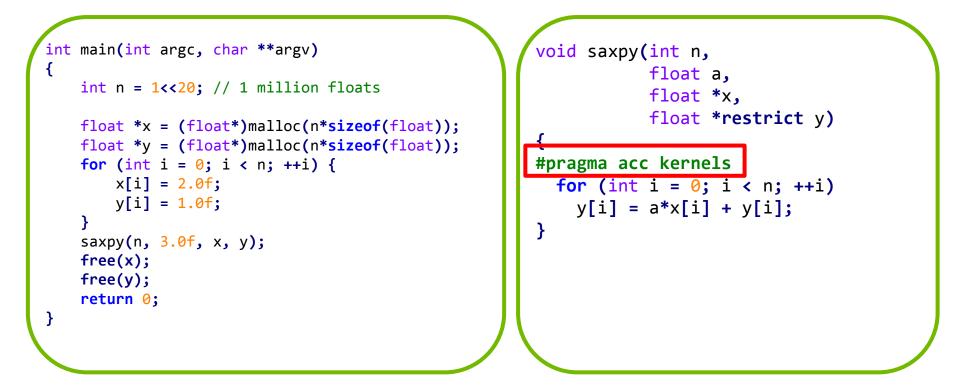






Complete saxpy.c

Only a single line to the above example is needed to produce an OpenACC SAXPY in C.









SAXPY code (only functions) in CUDA C

```
// define CUDA kernel function
global void saxpy kernel( float a, float* x, float* y, int n ){
    int i;
    i = blockIdx.x*blockDim.x + threadIdx.x;
    if( i <= n ) y[i] = a*x[i] + y[i];
}
void saxpy( float a, float* x, float* y, int n ){
    float *xd, *yd;
    // manage device memory
    cudaMalloc( (void**)&xd, n*sizeof(float) );
    cudaMalloc( (void**)&yd, n*sizeof(float) );
    cudaMemcpy( xd, x, n*sizeof(float), cudaMemcpyHostToDevice );
    cudaMemcpy( yd, y, n*sizeof(float), cudaMemcpyHostToDevice );
    // calls the kernel function
    saxpy_kernel<<< (n+31)/32, 32 >>>( a, xd, yd, n );
    cudaMemcpy( x, xd, n*sizeof(float), cudaMemcpyDeviceToHost );
    // free device memory after use
    cudaFree( xd );
    cudaFree( yd );
```

```
}
```





CUDA C/OpenACC – Big Difference

- With CUDA, we changed the structure of the old code. Non-CUDA programmers can't understand new code. It is not even ANSI standard code.
 - We have separate sections for the host code, and the GPU device code. Different flow of code. Serial path now gone forever.
 - Although CUDA C gives you maximum flexibility, the effort needed for restructuring the code seems to be high.
 - OpenACC seems ideal for researchers in science and engineering.





≻ C



Compiler output of the first example

Emit information about accelerator region targeting.

pgcc -acc -Minfo=accel -ta=nvidia,time saxpy_1stexample.c

> Fortran

pgf90 -acc -Minfo=accel -ta=nvidia,time saxpy_1stexample.c

- > Use "man pgcc/pgf90" to check the meaning of the compiler switches.
- Compiler output :

```
pgcc -acc -Minfo=accel -ta=nvidia,time saxpy_1stexample.c
saxpy:
    26, Generating present_or_copyin(x[:n])
        Generating present_or_copy(y[:n])
        Generating NVIDIA code
    27, Loop is parallelizable
        Accelerator kernel generated
        27, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */
```

Introduction to GPU Programming





Add PGI compiler to your environment

```
[fchen14@mike424 gpuex]$ cat ~/.soft
# This is the .soft file.
# It is used to customize your environment by setting up environment
# variables such as PATH and MANPATH.
# To learn what can be in this file, use 'man softenv'.
+portland-14.3
@default
[fchen14@mike424 gpuex]$ resoft
[fchen14@mike424 gpuex]$ pgcc -V
[fchen14@mike424 gpuex]$ cp -r /home/fchen14/gpuex/ ./
pgcc 14.3-0 64-bit target on x86-64 Linux -tp sandybridge
The Portland Group - PGI Compilers and Tools
Copyright (c) 2014, NVIDIA CORPORATION. All rights reserved.
[fchen14@mike424 gpuex]$ cd ~/gpuex
[fchen14@mike424 gpuex]$ cat saxpy 1stexample.c
[fchen14@mike424 gpuex]$ pgcc -acc -Minfo=accel -ta=nvidia,time
saxpy 1stexample.c
```



[fchen14@mike424 gpuex]\$./a.out



Runtime output

Accelerator Kernel Timing data /home/fchen14/loniworkshop2014/laplace/openacc/c/saxpy_1stexample.c saxpy NVIDIA devicenum=0 time(us): 2,247____ 2,247 = 1,421 + 637 + 18926: data region reached 1 time 26: data copyin reached 2 times device time(us): total=1,421 max=720 min=791 avg=710 29: data copyout reached 1 time device time(us): total=637 max=637 min=637 avg=637 26: compute region reached 1 time 26: kernel launched 1 time grid: [4096] block: [256] device time(us): total=189 max=189 min=189 avg=189 elapsed time(us): total=201 max=201 min=201 avg=201





OpenACC kernels directive

What is a kernel? A function that runs in parallel on the GPU.

- The kernels directive expresses that a region may contain parallelism and the compiler determines what can be safely parallelized.
- The compiler breaks code in the kernel region into a sequence of kernels for execution on the accelerator device.
- When a program encounters a kernels construct, it will launch a sequence of kernels in order on the device.
- The compiler identifies 2 parallel loops and generates 2 kernels below.

#pragma	acc kernels
for	(i = 0; i < n; i++){
	x[i] = 1.0;
	y[i] = 2.0;
}	
for	(i = 0; i < n; i++){
	y[i] = a*x[i] + y[i];
}	
}	





OpenACC parallel directive

- Similar to OpenMP, the parallel directive identifies a block of code as having parallelism.
- > Compiler generates one parallel kernel for that loop.
- ≻ C

```
#pragma acc parallel [clauses]
```

> Fortran

```
!$acc parallel [clauses]
```

```
#pragma acc parallel
{
    for (i = 0; i < n; i++){
        x[i] = 1.0;
        y[i] = 2.0;
    }
    for (i = 0; i < n; i++){
        y[i] = a*x[i] + y[i];
    }
}</pre>
```

```
!$acc parallel
do i = 1, n
    x(i) = 1.0
    y(i) = 2.0
end do
do i = 1, n
    y(i) = y(i) + a * x(i)
end do
!$acc end parallel
```







OpenACC loop directive

Loops are the most likely targets for parallelizing.

- The Loop directive is used within a parallel or kernels directive identifying a loop that can be executed on the accelerator device.
- The loop directive can be combined with the enclosing parallel or kernels
- The loop directive clauses can be used to optimize the code. This however requires knowledge of the accelerator device.
- Clauses: gang, worker, vector, num_gangs, num_workers
- > C: #pragma acc [parallel/kernels] loop [clauses]
- Fortran: !\$acc [parallel/kernels] loop [clauses]

```
#pragma acc loop
for (i = 0; i < n; i++){
    y[i] = a*x[i] + y[i];
}</pre>
```

```
!$acc loop
do i = 1, n
    y(i) = y(i) + a * x(i)
end do
!$acc end loop
```





OpenACC kernels vs parallel

> kernels

- Compiler performs parallel analysis and parallelizes what it believes is safe.
- Can cover larger area of code with single directive.
- > parallel
 - Requires analysis by programmer to ensure safe parallelism.
 - Straightforward path from OpenMP
- Both approaches are equally valid and can perform equally well.







Clauses

- data management clauses
 - copy(...),copyin(...), copyout(...)
 - create(...), present(...)
 - present_or_copy{,in,out}(...) or pcopy{,in,out}(...)
 - present_or_create(...) or pcreate(...)
- > reduction(operator:list)
- > if (condition)
- > async (expression)







Runtime Libraries

System setup routines

- acc_init(acc_device_nvidia)
- acc_set_device_type(acc_device_nvidia)
- acc_set_device_num(acc_device_nvidia)
- Synchronization routines
 - acc_async_wait(int)
 - acc_async_wait_all()

For more information, refer to the OpenACC standard







Second example: Jacobi Iteration

Solve Laplace equation in 2D:

 Iteratively converges to correct value (e.g. Temperature), by computing new values at each point from the average of neighboring points.

$$\nabla^2 f(x,y) = 0$$

$$A_{k+1}(i,j) = \frac{A_k(i-1,j) + A_k(i+1,j) + A_k(i,j-1) + A_k(i,j+1)}{4}$$

$$A(i,j+1)$$

$$A(i-1,j)$$

$$A(i,j-1)$$

$$A(i+1,j)$$







Graphical representation for Jacobi iteration

Current Array: A

1.0	1.0	1.0	1.0	1.0	1.0	1.0	1.0
1.0	2.0	4.0	6.0	8.0	10.0	12.0	1.0
1.0	3.0	5.0	7.0	9.0	11.0	13.0	1.0
1.0	2.0	6.0	1.0	3.0	7.0	5.0	1.0
1.0	1.0	1.0	1.0	1.0	1.0	1.0	1.0

Next Array: Anew

1.0	1.0	1.0	1.0	1.0	1.0	1.0	1.0
1.0	2.25	3.56	6.0				1.0
1.0		5.0					1.0
1.0							1.0
1.0	1.0	1.0	1.0	1.0	1.0	1.0	1.0

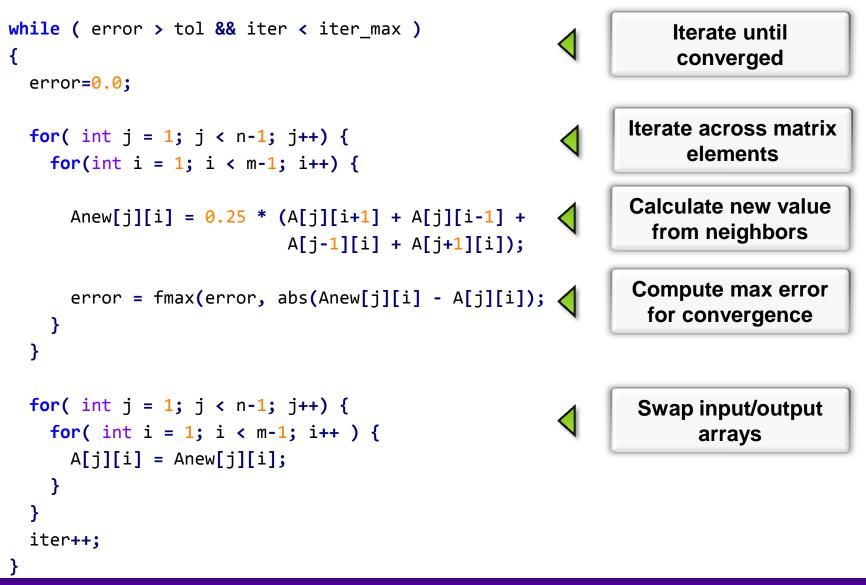




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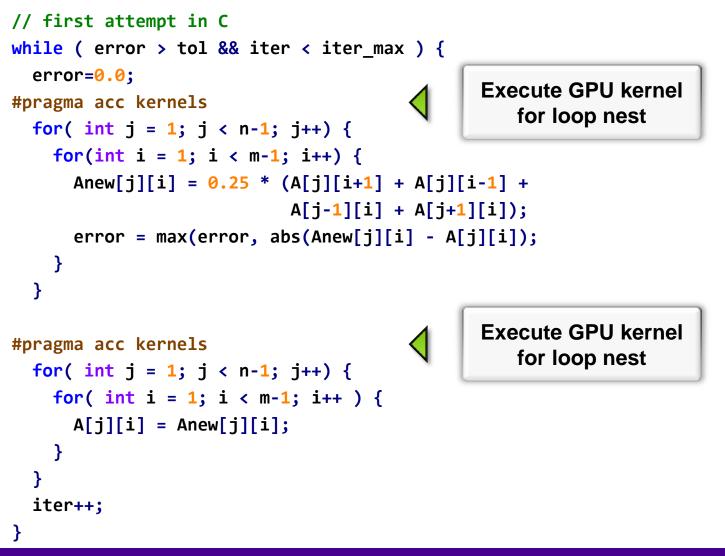
Serial version of the Jacobi Iteration







First Attempt in OpenACC







Compiler Output

pgcc -acc -Minfo=accel -ta=nvidia,time laplace_openacc.c -o laplace_acc.out
main:

65, Generating present or copyin(Anew[1:4094][1:4094]) Generating present or copyin(A[:4096][:4096]) present_or_copyin Generating NVIDIA code 66, Loop is parallelizable 67, Loop is parallelizable Accelerator kernel generated 66, #pragma acc loop gang /* blockIdx.y */ 67, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */ 70 Max reduction generated for error 75, Generating present or copyin(Anew[1:4094][1:4094]) Generating present or copyin(A[1:4094][1:4094]) present_or_copyin Generating NVIDIA code 76, Loop is parallelizable 77, Loop is parallelizable Accelerator kernel generated 76, #pragma acc loop gang /* blockIdx.y */ 77, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */





Performance of First Jacobi ACC Attempt

- > CPU: Intel(R) Xeon(R) CPU E5-2670 @ 2.60GHz
- GPU: Nvidia Tesla K20Xm
- The OpenACC code is even slower than the single thread/serial version of the code
- > What is the reason for the significant slow-down?

Execution	Time (sec)	Speedup
OpenMP 1 threads	45.64	
OpenMP 2 threads	30.05	1.52
OpenMP 4 threads	24.91	1.83
OpenMP 8 threads	25.24	1.81
OpenMP 16 threads	26.19	1.74
OpenACC w/GPU	190.32	0.24







Output Timing Information from Profiler

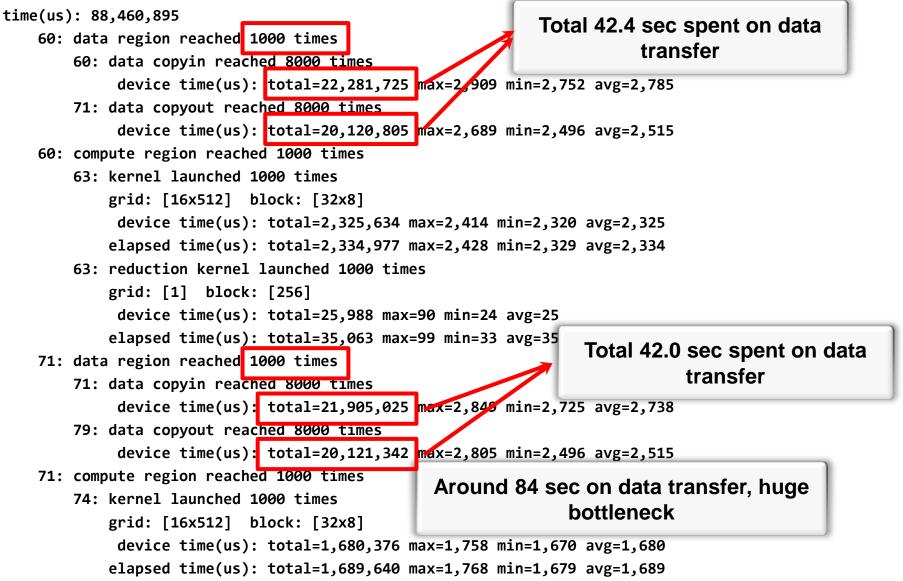
- > Use compiler flag: -ta=nvidia, time
 - Link with a profile library to collect simple timing information for accelerator regions.
- OR set environmental variable: export PGI_ACC_TIME=1
 - Enables the same lightweight profiler to measure data movement and accelerator kernel execution time and print a summary at the end of program execution.
- Either way can output profiling information







Accelerator Kernel Timing data (1st attempt)





Introduction to GPU Programming

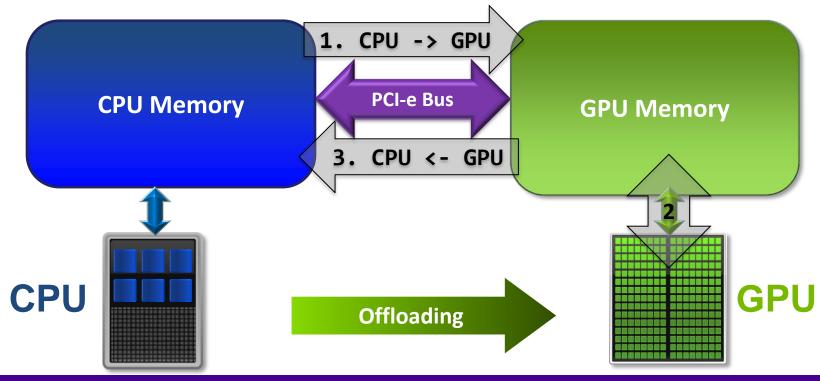


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Recall Basic Concepts on Offloading

- CPU and GPU have their respective memory, connected through PCI-e bus
- Processing Flow of the offloading
 - 1. Copy input data from CPU memory to GPU memory
 - 2. Load GPU program and execute
 - 3. Copy results from GPU memory to CPU memory

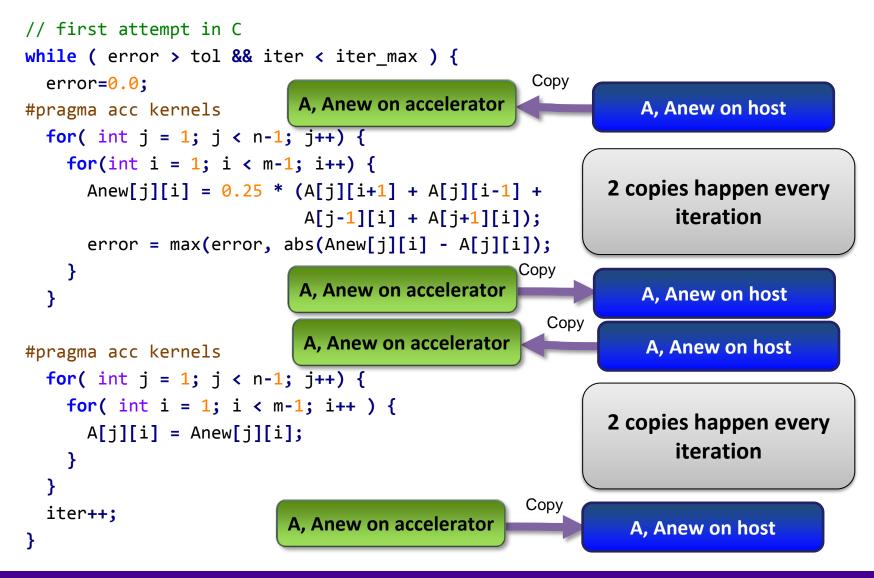




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Excessive Data Transfers







Rules of Coprocessor (GPU) Programming

- Transfer the data across the PCI-e bus onto the device and keep it there.
- Give the device enough work to do (avoid preparing data).
- Focus on data reuse within the coprocessor(s) to avoid memory bandwidth bottlenecks.







OpenACC Data Management with Data Region

C syntax

- #pragma acc data [clause]
- { structured block/statement }

Fortran syntax

- !\$acc data [clause]
- structured block
- !\$acc end data
- Data regions may be nested.







Data Clauses

> copy (list)

/* Allocates memory on GPU and copies data from host to GPU
when entering region and copies data to the host when exiting region.*/

> copyin (list)

/* Allocates memory on GPU and copies data from host to GPU when entering region. */

> copyout (list)

/* Allocates memory on GPU and copies data to the host when exiting region. */

> create (list)

/* Allocates memory on GPU but does not copy. */

> present (list)

/* Data is already present on GPU from another containing data region.
*/

> and present_or_copy[in|out], present_or_create, deviceptr.





Second Attempt: OpenACC C

```
#pragma acc data copy(A), create(Anew)
while ( error > tol && iter < iter_max ) {
    error=0.0;</pre>
```

Copy A in at beginning of loop, out at end. Allocate Anew on accelerator

```
#pragma acc kernels
  for( int j = 1; j < n-1; j++) {</pre>
    for(int i = 1; i < m-1; i++) {</pre>
      Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                             A[j-1][i] + A[j+1][i]);
      error = max(error, abs(Anew[j][i] - A[j][i]);
    }
  }
#pragma acc kernels
  for( int j = 1; j < n-1; j++) {</pre>
    for( int i = 1; i < m-1; i++ ) {</pre>
      A[j][i] = Anew[j][i];
    }
  }
  iter++;
```

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Second Attempt: OpenACC Fortran

```
Copy A in at beginning of loop,
 !$acc data copy(A), create(Anew)
                                                                                                                                                                                                                                                                                                                                   out at end. Allocate Anew on
do while ( err > tol .and. iter < iter max</pre>
                                                                                                                                                                                                                                                                                                                                                                                 accelerator
             err=0._fp_kind
 !$acc kernels
             do j=1,m
                          do i=1,n
                                       Anew(i,j) = .25_{fp_kind} * (A(i+1, j) + A(i-1, j) + & (A(i+1, j)) + & (A(i+
                                                                                                                                                                                                                        A(i, j-1) + A(i, j+1))
                                       err = max(err, Anew(i,j) - A(i,j))
                          end do
             end do
 !$acc end kernels
  . . .
iter = iter +1
end do
 !$acc end data
```





Second Attempt: Performance

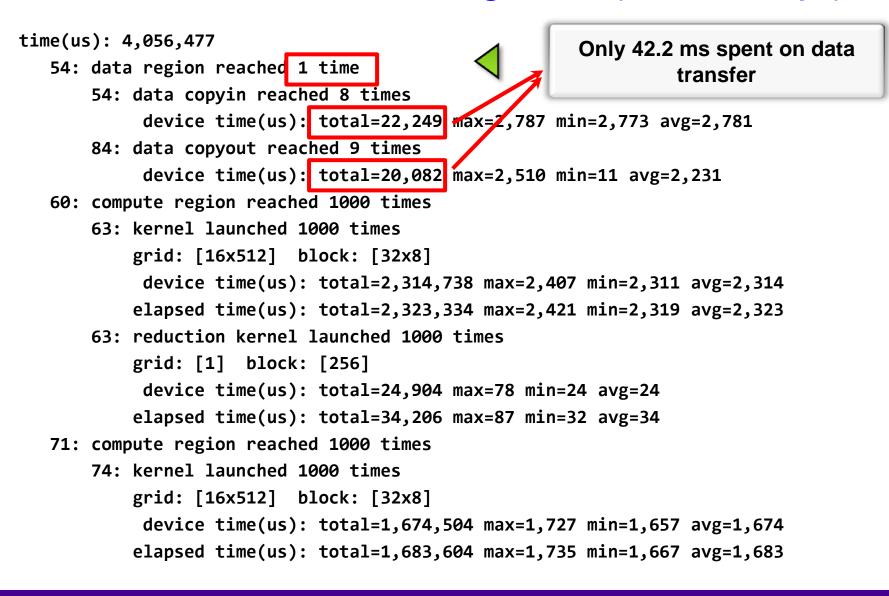
- Significant speedup after the insertion of the data region directive
- > CPU: Intel Xeon CPU E5-2670 @ 2.60GHz
- GPU: Nvidia Tesla K20Xm

Execution	Time (sec)	Speedup
OpenMP 1 threads	45.64	
OpenMP 2 threads	30.05	1.52
OpenMP 4 threads	24.91	1.83
OpenACC w/GPU (data region)	4.47	10.21 (serial) 5.57 (4 threads)





Accelerator Kernel Timing data (2nd attempt)







Array Shaping

- Compiler sometimes cannot determine size of arrays
 - Sometimes we just need to use a portion of the arrays
 - we will see this example in the exercise
- Under such case, we must specify explicitly using data clauses and array "shape" for this case
- ≻ C

```
#pragma acc data copyin(a[0:size]), copyout(b[s/4:3*s/4])
```

Fortran

```
!$pragma acc data copyin(a(1:size)), copyout(b(s/4:3*s/4))
```

- The number between brackets are the beginning element followed by the number of elements to copy:
 - [start_element:number_of_elements_to_copy]
 - In C/C++, this means start at a[0] and continue for "size" elements.
- Note: data clauses can be used on data, kernels or parallel





Update Construct

> Fortran

```
#pragma acc update [clause ...]
```

```
≻ C
```

```
!$acc update [clause ...]
```

- Used to update existing data after it has changed in its corresponding copy (e.g. update device copy after host copy changes)
- Move data from GPU to host, or host to GPU. Data movement can be conditional, and asynchronous.







Further Speedups

- OpenACC gives us more detailed control over parallelization via gang, worker, and vector clauses
 - PE (processing element) as a SM (streaming multiprocessor)
 - gang == CUDA threadblock
 - worker == CUDA warp
 - vector == CUDA thread
- By understanding more about OpenACC execution model and GPU hardware organization, we can get higher speedups on this code
- By understanding bottlenecks in the code via profiling, we can reorganize the code for higher performance







Finding Parallelism in your code

- (Nested) for loops are best for parallelization
 - Large loop counts needed to offset GPU/memcpy overhead
- Iterations of loops must be independent of each other
 - To help compiler:
 - restrict keyword
 - independent clause
- Compiler must be able to figure out sizes of data regions
 - Can use directives to explicitly control sizes
- Pointer arithmetic should be avoided if possible
 - Use subscripted arrays, rather than pointer-indexed arrays.
- Function calls within accelerated region must be inlineable.







Exercise 1

For the matrix multiplication code $A \cdot B = C$

where:

$$a_{i,j} = i + j$$

$$b_{i,j} = i \cdot j$$

$$c_{i,j} = \sum_{k} a_{i,k} \cdot b_{k,j}$$

- 1. For mm_acc_v0.c, speedup the matrix multiplication code segment using OpenACC directives
- 2. For mm_acc_v1.c:
 - Change A, B and C to dynamic arrays, i.e., the size of the matrix can be specified at runtime;
 - Complete the function matmul_acc using the OpenACC directives;
 - Compare performance with serial and OpenMP results





Exercise 2

- > Complete the saxpy example using OpenACC directives. $\vec{y} = a \cdot \vec{x} + \vec{y}$
- Calculate the result of a constant times a vector plus a vector:
 - where *a* is a constant, \vec{x} and \vec{y} are one dimensional vectors.
 - 1. Add OpenACC directives for initialization of x and y arrays;
 - 2. Add OpenACC directives for the code for the vector addition;
 - 3. Compare the performance with OpenMP results;







Exercise 3

> Calculate π value using the equation:

$$\int_{0}^{1} \frac{4.0}{(1.0 + x^2)} = \pi$$

with the numerical integration:
$$\sum_{i=1}^{n} \frac{4.0}{(1.0 + x_i \cdot x_i)} \Delta x \approx \pi$$

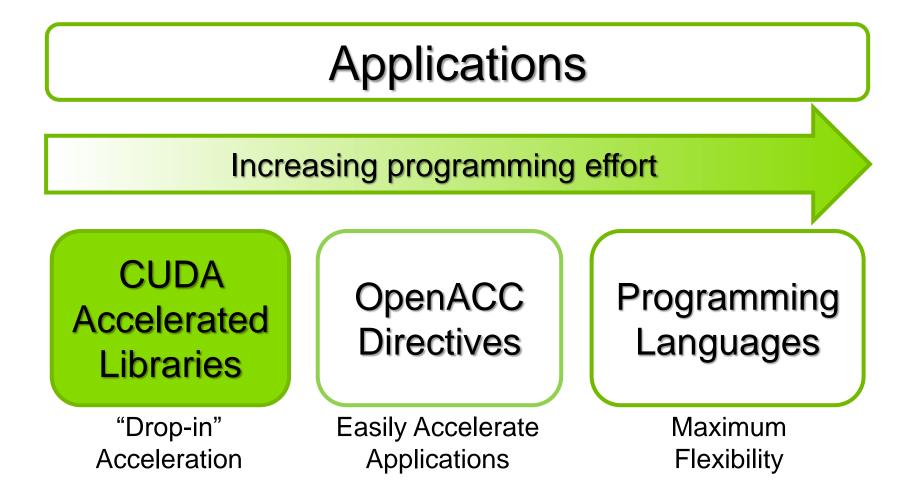
1. Complete the code using OpenACC directives







3 Ways to Accelerate Applications







Drop-In Acceleration (Step 1)

int N = 1 << 20;

// Perform SAXPY on 1M elements: d_y[]=a*d_x[]+d_y[]

cublasSaxpy(h, N, &alpha, d_x, 1, d_y, 1);



Add "cublas" prefix and use device variables







Drop-In Acceleration (Step 2)

int N = 1 << 20; cublasHandle_t h; cublasCreate(&h);

 \triangleleft

Initialize CUBLAS

// Perform SAXPY on 1M elements: d_y[]=a*d_x[]+d_y[]
cublasSaxpy(h, N, &alpha, d_x, 1, d_y, 1);

cublasDestroy(h); cudaDeviceReset();



Shut down CUBLAS







Drop-In Acceleration (Step 3)

```
int N = 1 << 20;
cublasHandle_t h;
cublasCreate(&h);
cudaMalloc((void**)&d x, N*sizeof(float));
cudaMalloc((void**)&d y, N*sizeof(float));
```



// Perform SAXPY on 1M elements: d_y[]=a*d_x[]+d_y[] cublasSaxpy(h, N, &alpha, d_x, 1, d_y, 1);

```
cudaFree(d_x);
cudaFree(d y);
cublasDestroy(h);
cudaDeviceReset();
```



Deallocate device vectors







Drop-In Acceleration (Step 4)

```
int N = 1 << 20;
cublasHandle t h;
cublasCreate(&h);
cudaMalloc((void**)&d_x, N*sizeof(float));
cudaMalloc((void**)&d y, N*sizeof(float));
cudaMemcpy(d_x, &x[0], N*sizeof(float), cudaMemcpyHostToDevice);
                                                                       Transfer
cudaMemcpy(d_y, &y[0], N*sizeof(float), cudaMemcpyHostToDevice);
                                                                      data to GPU
// Perform SAXPY on 1M elements: d_y[]=a*d_x[]+d_y[]
cublasSaxpy(h, N, &alpha, d_x, 1, d_y, 1);
                                                                       Read data
cudaMemcpy(&y[0], d_y, N*sizeof(float), cudaMemcpyDeviceToHost);
                                                                       back GPU
cudaFree(d x);
cudaFree(d_y);
cublasDestroy(h);
cudaDeviceReset();
```





Compile and Run

Need to link to the cublas library

[fchen14@mike424 gpuex]\$ nvcc cublas_vec_add.cu -1 cublas
[fchen14@mike424 gpuex]\$

> Run example:

[fchen14@mike424 gpuex]\$./a.out
cublas time took 0.307 ms
x[0] = 7.200000
y[0] = 5.300000
z[0] = 12.500000

