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Introduction to OpenMP

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- Acquaint users with the concept of shared memory parallelism
- Acquaint users with the basics of programming with OpenMP

Goals







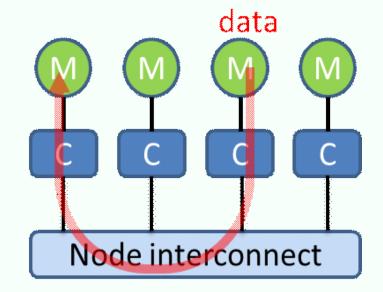
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- Each process has its own address space
 - Data is local to each process
- Data sharing achieved via explicit message passing
- Example

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• MPI









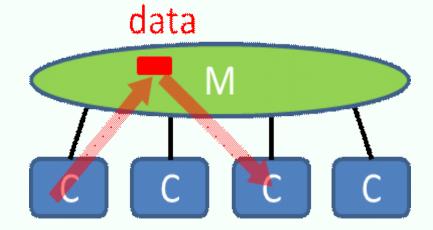
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Shared memory model

- All threads can access the global memory space
- Data sharing achieved via writing to/reading from the same memory location
- Example

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- OpenMP
- Pthreads







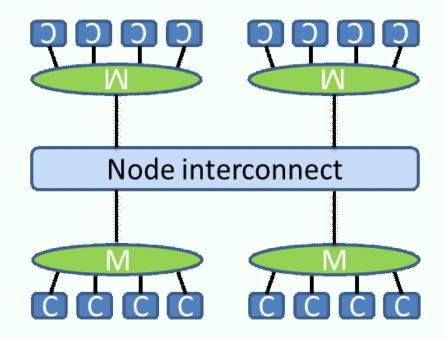




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Clusters of SMP nodes

- The shared memory model is most commonly represented by Symmetric Multi-Processing (SMP) systems
 - Identical processors
 - Equal access time to memory
- Large shared memory systems are rare, clusters of SMP nodes are popular









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Shared vs Distributed

Distributed Memory

Shared Memory

• Pros

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- Memory scalable with number of number of processors
- Easier and cheaper to build
- Cons
 - Difficult load balancing
 - Data sharing is slow

• Pros

- Global address space is userfriendly
- Data sharing is fast
- Cons
 - Lack of scalability
 - Data conflict issues









OpenMP

- OpenMP is an Application Program Interface (API) for thread based parallelism; Supports Fortran, C and C++
- Uses a fork-join execution model
- OpenMP structures are built with program directives, runtime libraries and environment variables
- OpenMP has been the industry standard for shared memory programming over the last decade
 - Permanent members of the OpenMP Architecture Review Board: AMD, Cray, Fujutsu, HP, IBM, Intel, Microsoft, NEC, PGI, SGI, Sun
- OpenMP 3.1 was released in September 2011











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Advantages of OpenMP

- Portability
 - Standard among many shared memory platforms
 - Implemented in major compiler suites
- Ease to use
 - Serial programs can be parallelized by adding compiler directives
 - Allows for incremental parallelization a serial program evolves into a parallel program by parallelizing different sections incrementally





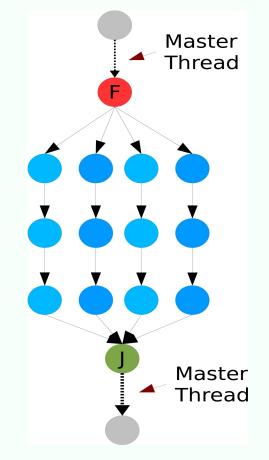


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Fork-join Execution Model

- Parallelism is achieved by generating multiple threads that run in parallel
 - A fork is when a single thread is made into multiple, concurrently executing threads
 - A join is when the concurrently executing threads synchronize back into a single thread
- OpenMP programs essentially consist of a series of forks and joins.







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Building Blocks of OpenMP

- Program directives
 - Syntax
 - C/C++: #pragma omp <directive> [clause]
 - Fortran: !\$omp <directive> [clause]
 - Parallel regions
 - Parallel loops
 - Synchronization
 - Data structure
 - .
- Runtime library routines
- Environment variables



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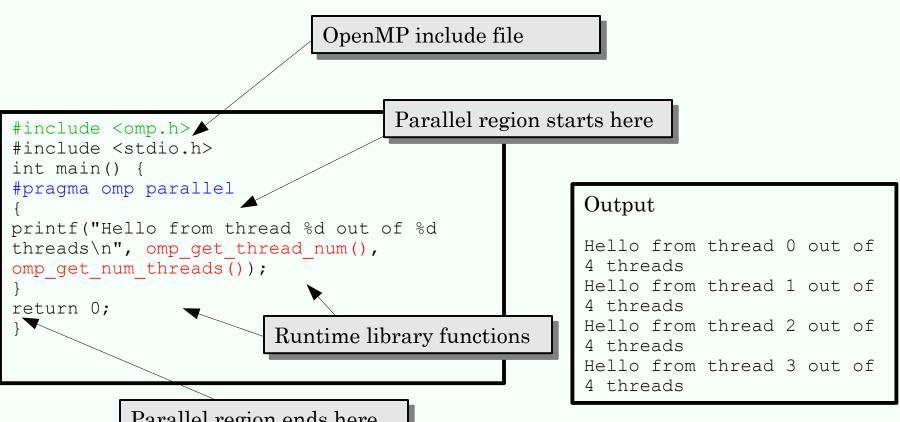




Hello World: C



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Parallel region ends here



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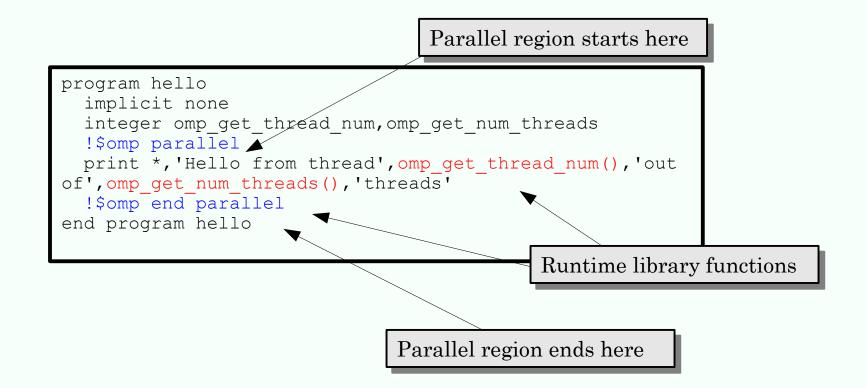






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Hello World: Fortran











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Compilation and Execution

- IBM p575 clusters
 - Use the thread-safe compilers (with "_r")
 - Use '-qsmp=omp' option

%xlc_r -qsmp=omp test.c && OMP_NUM_THREADS=4 ./a.out

• Dell Linux clusters

• Use '-openmp' option (Intel compiler)

%icc -openmp test.c && OMP_NUM_THREADS=4 ./a.out











Exercise 1: Hello World

- Write a "hello world" program with OpenMP where
 - If the thread id is odd, then print a message "Hello world from thread x, I'm odd!"
 - If the thread id is even, then print a message "Hello world from thread x, I am even!"







Solution



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C/C++ Fortran #include <omp.h> program hello #include <stdio.h> implicit none int main() { integer i, omp get thread num int id; !\$omp parallel private(i) #pragma omp parallel private(id) i = omp get thread num() if (mod(i,2).eq.1) then print *, 'Hello world from id = omp get thread num(); if (id%2==1) thread', i, ', I am odd!' printf("Hello world from else thread %d, I am odd\n", id); print *, 'Hello world from thread', i, ', I am even!' else printf("Hello world from endif thread %d, I am evenn, id); !\$omp end parallel end program hello









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Work Sharing: Parallel Loops

- We need to share work among threads to achieve parallelism
- Loops are the most likely targets when parallelizing a serial program
- Syntax

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- Fortran: !\$omp parallel do
- C/C++: #pragma omp parallel for
- Other working sharing directives available
 - Sections (discussed later)
 - Tasks (new feature in OpenMP 3.0)



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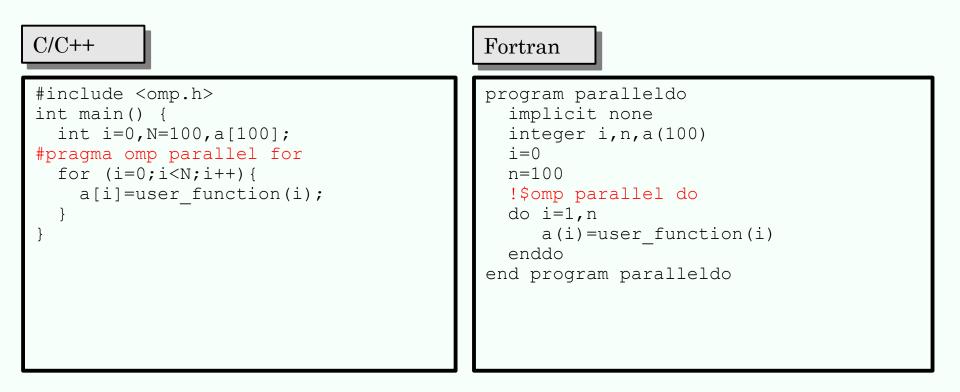




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Example: Parallel Loops











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Load Balancing (1)

- OpenMP provides different methods to divide iterations among threads, indicated by the schedule clause
 - Syntax: schedule (<method>, [chunk size])
- Methods include
 - Static: the default schedule; divide iterations into chunks according to size, then distribute chunks to each thread in a round-robin manner;
 - Dynamic: each thread grabs a chunk of iterations, then requests another chunk upon the completion of the current one, until all iterations executed
 - Guided: similar to dynamic; the only difference is that the chunk size starts large and shrinks to size eventually



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Load Balancing (2)

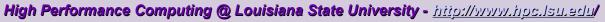
4 threads, 100 iterations

Schedule	Iterations mapped onto thread				
	0	1	2	3	
Static	1-25	26-50	51-75	76-100	-
Static,20	1-20,81-100	21-40	41-60	61-80	
Dynamic	1	2	3	4	
Dynamic,10	1-10	11-20	21-30	31-40	



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Load Balancing (3)

Schedule	When to use	
Static	Even and predictable workload per iteration; scheduling may be done at compilation time, least work at runtime;	
Dynamic	Highly variable and unpredictable workload per iteration; most work at runtime	
Guided	Special case of dynamic scheduling; compromise between load balancing and scheduling overhead at runtime	









Working Sharing: Sections

• Gives a different block to each thread

C/C++	Fortran
<pre>#pragma omp parallel { #pragma omp sections { #pragma omp section some_calculation(); #pragma omp section more_calculation(); #pragma omp section yet_more_calculation(); } }</pre>	<pre>!\$omp parallel !\$omp sections !\$omp section call some_calculation !\$omp section call more_calculation !\$omp section call yet_more_calculation !\$omp end sections !\$omp end parallel</pre>





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Scope of Variables

- Shared(list)
 - Specifies the variables that are shared among all the threads
- Private(list)
 - Creates a local copy of the specified variables for each thread
 - the value is uninitialized!
- Default(shared|private|none)
 - Defines the default scope of variables
 - C/C++ API does not have default (private)
- Most variables are shared by default
 - A few exceptions: iteration variables; stack variables in subroutines; automatic variables within a statement block









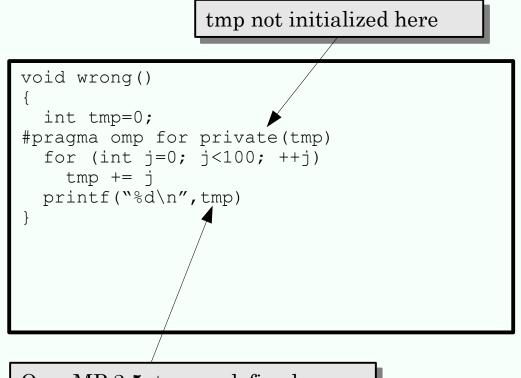
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Private Variables

- Not initialized at the beginning of parallel region
- After the parallel region

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- Not defined in OpenMP 2.x
- 0 in OpenMP 3.x



OpenMP 2.5: tmp undefined OpenMP 3.0: tmp is 0

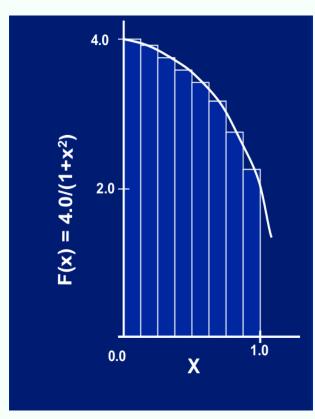






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We know that:

$$\int_{0}^{1} \frac{4.0}{(1+x^2)} \, dx = \pi$$

So numerically we can approximate pi as the sum of the area of a number of rectangles



Source: Meadows et al, A "hands-on" introduction to OpenMP, SC09



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Exercise 2: serial version

C/C++	Fortran
<pre>double x,deltax,pi,sum=0.0 int i,nstep=</pre>	Real*8 :: x,deltax,pi,sum integer :: i,nstep
<pre>deltax=1./(double)nstep</pre>	nstep= sum=0
<pre>for (i=0; i<nstep; i++)="" x="(i+0.5)*deltax</pre" {=""></nstep;></pre>	<pre>deltax=1./float(nstep)</pre>
<pre>sum=sum+4./(1.+x*x) }</pre>	<pre>do i=1,nstep x=(i+0.5)*deltax sum=sum+4./(1.+x*x)</pre>
pi=deltax*sum	enddo pi=deltax*sum









Exercise 2: OpenMP version

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- Create a parallel version of the program with OpenMP







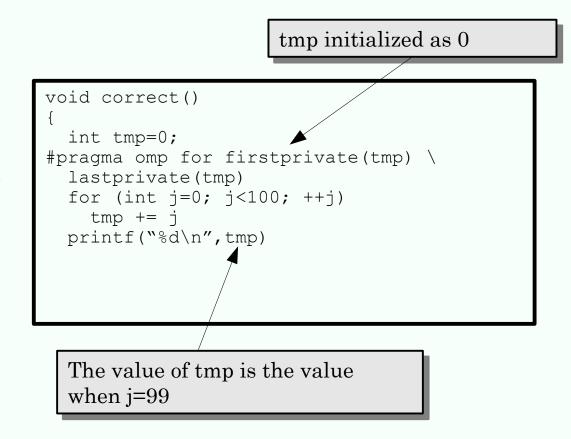


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Special Cases of Private

- Firstprivate
 - Initialize each private copy with the corresponding value from the master thread
- Lastprivate
 - Allows the value of a private variable to be passed to the shared variable outside the parallel region











Reduction

- The reduction clause allows accumulative operations on the value of variables
- Syntax: reduction (operator: variable list)
- A private copy of each variable appears in reduction is created as if the private clause is specified
- Operators
 - Arithmetic
 - Bitwise
 - Logical









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Example: Reduction

C/C++	Fortran
<pre>#include <omp.h></omp.h></pre>	<pre>program reduction</pre>
int main() {	implicit none
int i,N=100,sum,a[100],b[100];	integer i,n,sum,a(100),b(100)
for (i=0;i <n;i++) td="" {<=""><td>n=100</td></n;i++)>	n=100
a[i]=i;	do i=1,n
b[i]=1;	a(i)=i
}	enddo
sum=0;	b=1
#pragma omp parallel for	sum=0
reduction(+:sum)	!\$omp parallel do reduction(+:sum)
for (i=0;i <n;i++) td="" {<=""><td>do i=1,n</td></n;i++)>	do i=1,n
sum=sum+a[i]*b[i];	sum=sum+a(i)*b(i)
}	enddo
}	end







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Exercise 3: pi calculation with reduction

• Redo exercise 2 with reduction



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Pitfalls (1): False Sharing

- Array elements that are in the same cache line can lead to false sharing
 - The system handles cache coherence on a cache line basis, not on a byte or word basis
 - Each update of a single element could invalidate the entire cache line









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Pitfalls (2): Race Condition

- Multiple threads try to write to the same memory location at the same time
 - Indeterministic results
- Inappropriate scope of variable can cause indeterministic results too
- When having indeterministic results, set the number of threads to 1 to check
 - If problem persists: scope problem
 - If problem is solved: race condition



!\$omp parallel do do i=1,n if (a(i).gt.max) then max=a(i) endif enddo







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Synchronization: Barrier

- "Stop sign" where every thread waits until all threads arrive
- Purpose: protect access to shared data
- Syntax
 - Fortran: !\$omp barrier
 - C/C++: #pragma omp barrier
- A barrier is implied at the end of every parallel region
 - Use the nowait clause to turn if off
- Synchronizations are costly so their usage should be minimized









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• Critical

• Only one thread at a time can enter a critical region

• Atomic

• Only one thread at a time can update a memory location

```
double x;
#pragma omp parallel for
for (i=0;i<N;i++)
{
    a = some_calculation(i)
#pragma omp critical
    some function(a,x);</pre>
```

```
double a;
#pragma omp parallel
{ double b;
   b = some_calculation();
#pragma omp atomic
   a += b;
}
```









Runtime Library Functions

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- Modify/query the number of threads
 - omp_set_num_threads(), omp_get_num_threads(), omp_get_thread_num(), omp_get_max_threads()
- Query the number of processes
 - omp_num_procs()
- Query whether or not in an active parallel region
 - omp_in_parallel()
- Control the behavior of dynamic threads
 - omp_set_dynamic(), omp_get_dynamic()



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Environment Variables

- OMP_NUM_THREADS: set default number of threads to use
- OMP_SCHEDULE: control how iterations are scheduled for parallel loops









References

- https://docs.loni.org/wiki/Using_OpenMP
- http://en.wikipedia.org/wiki/OpenMP
- http://www.nersc.gov/nusers/help/tutorials/openmp/
- http://www.llnl.gov/computing/tutorials/openMP/









Next week's training

- What: Shell scripting tutorial
- Where: Frey 307
- When: Oct 26, Wednesday, 10am-12pm



