HPC User Environment 2

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Outline

➢ Review HPC User Environment 1 topics
   – Available HPC resources
   – Accounts and Allocations
   – Cluster architecture
   – Connect to clusters
   – Software management using softenv and module

➢ Things to be covered in this training
   – Job management
     • Interactive vs Batch jobs
     • Submit and monitor your jobs
   – Understanding Job scheduling
     • Job priority
     • Backfill
   – Compiling and analyze codes on cluster
     • Serial program
     • Parallel program
Brief Review of Session 1
Inside A Cluster Rack

- Rack
- Infiniband Switch
- Compute Node
Inside A QB2 Compute Node (Dell C8000)

- Accelerator 1 (GPU)
- Accelerator 2 (GPU)
- Storage
- Processor
- Processor
- Network Card
- Memory
Conceptual Relationship

Cluster

Compute Node

CPUs

GPUs

Memory

LONI User

Head Node

Computation Nodes

File Server

Storage

Switch

Switch

Internet
## Cluster Nomenclature

<table>
<thead>
<tr>
<th>Term</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Cluster</td>
<td>The top-level organizational unit of an HPC cluster, comprising a set of nodes, a queue, and jobs.</td>
</tr>
<tr>
<td>Node</td>
<td>A single, named host machine in the cluster.</td>
</tr>
<tr>
<td>Core</td>
<td>The basic computation unit of the CPU. For example, a quad-core processor is considered 4 cores.</td>
</tr>
<tr>
<td>Job</td>
<td>A user's request to use a certain amount of resources for a certain amount of time on cluster for his work.</td>
</tr>
</tbody>
</table>
HPC Cluster Architectures

➢ Major architecture
   – Intel x86_64 clusters
     • Vendor: Dell
     • Operating System: Linux (RHEL 4/5/6)
     • Processor: Intel
Accessing cluster using ssh (Secure Shell)

➢ **On Unix and Mac**
  – use ssh on a terminal to connect

➢ **Windows box (ssh client):**
  – Putty, Cygwin ([http://www.chiark.greenend.org.uk/~sgtatham/putty/download.html](http://www.chiark.greenend.org.uk/~sgtatham/putty/download.html))

➢ **ssh username@mike.hpc.lsu.edu**

➢ **Host name**
  – LONI: `<cluster_name>.loni.org`
    • `<cluster_name>` can be:
      – `qb.loni.org`
  – LSU HPC: `<cluster_name>.hpc.lsu.edu`
    • `<cluster_name>` can be:
      – `mike.hpc.lsu.edu`
      – `smic.hpc.lsu.edu`
      – `philip.hpc.lsu.edu`
Review Questions for Section 1
Access to cluster

How do I connect to HPC/LONI cluster?

a) By logging onto HPC webpage at www.hpc.lsu.edu
b) Using an ssh (secure shell) client such as MobaXterm/Putty
c) Go to the machine room in ISB in downtown Baton Rouge and connect my laptop to the nodes using a cable 😊
Software Management

➢ How do we manage the software installed on HPC/LONI clusters?
  – Using the modules command
  – Using a drop down menu on the www.hpc.lsu.edu webpage

➢ Recall the basic 5 module commands, what are they used for?
  – module av
  – module li
  – module disp
  – module load/unload <key>
  – module swap <key1> <key2>
Account and Allocation Policy

➢ Who can apply for allocations?
   a) Graduate student
   b) PostDoc
   c) Full time faculty
   d) All of the above
Account Security

➢ How to get your account suspended?
   a) Give your password to your friend/lab colleagues and let him/her use it.
   b) Give your password to your advisor so he/she can use your account to see your data.
   c) Run my simulation on the login node.
   d) All of the above
HPC User Environment 2

More on Job Queues
Cluster Environment

- Multiple compute nodes
- Multiple users
- Each user may have multiple jobs running simultaneously
- Multiple users may share the same node
Job submission basics

1. Find appropriate queue

2. Understand the queuing system and your requirements and proceed to submit jobs

3. Monitor jobs during execution
Job Queues

➢ Nodes are organized into queues. Nodes can be shared.
➢ Each job queue differs in
  – Number of available nodes
  – Max run time
  – Max running jobs per user
  – Nodes may have special characteristics: GPU/Xeon Phi’s, Large memory, etc.
➢ Jobs need to specify resource requirements
  – Nodes, time, queue
➢ It’s called a queue for a reason, but jobs don’t run on a “First Come First Served” policy,
  – This will be detailed in later slides
## Queue Characteristics – LONI clusters

<table>
<thead>
<tr>
<th>Machine</th>
<th>Queue</th>
<th>Max Runtime</th>
<th>ppn</th>
<th>Max running jobs</th>
<th>Max nodes per job</th>
<th>Use</th>
</tr>
</thead>
<tbody>
<tr>
<td>QB2</td>
<td>workq</td>
<td>3 days</td>
<td>20</td>
<td>44</td>
<td>128</td>
<td>Unpreemptable</td>
</tr>
<tr>
<td></td>
<td>checkpoint</td>
<td>20</td>
<td></td>
<td></td>
<td>256</td>
<td>Preemptable</td>
</tr>
<tr>
<td></td>
<td>single</td>
<td>7 days</td>
<td>1,2,4,6,8</td>
<td>1</td>
<td>ppn=1/2/4/6/8</td>
<td></td>
</tr>
</tbody>
</table>
## Queue Characteristics – LSU Linux clusters

<table>
<thead>
<tr>
<th>Machine</th>
<th>Queue</th>
<th>Max Runtime</th>
<th>ppn</th>
<th>Max running jobs</th>
<th>Max nodes per job</th>
<th>Use</th>
</tr>
</thead>
<tbody>
<tr>
<td>SuperMike II</td>
<td>workq</td>
<td>3 days</td>
<td>16</td>
<td>34</td>
<td>128</td>
<td>Unpreemptable</td>
</tr>
<tr>
<td></td>
<td>checkpt</td>
<td></td>
<td>16</td>
<td></td>
<td>128</td>
<td>Preemptable</td>
</tr>
<tr>
<td></td>
<td>bigmem</td>
<td>2 days</td>
<td>16</td>
<td></td>
<td>1</td>
<td>Big memory</td>
</tr>
<tr>
<td></td>
<td>gpu</td>
<td>3 days</td>
<td>16</td>
<td></td>
<td>16</td>
<td>Job using GPU</td>
</tr>
<tr>
<td></td>
<td>single</td>
<td>3 days</td>
<td>1,2,4,8</td>
<td></td>
<td>1</td>
<td>Single node jobs</td>
</tr>
<tr>
<td>Philip</td>
<td>workq</td>
<td>3 days</td>
<td>8</td>
<td>5</td>
<td>4</td>
<td>Unpreemptable</td>
</tr>
<tr>
<td></td>
<td>checkpt</td>
<td></td>
<td>8</td>
<td></td>
<td>4</td>
<td>Preemptable</td>
</tr>
<tr>
<td></td>
<td>bigmem</td>
<td></td>
<td>8</td>
<td></td>
<td>2</td>
<td>Big memory</td>
</tr>
<tr>
<td></td>
<td>single</td>
<td>14 days</td>
<td>4</td>
<td>50</td>
<td>1</td>
<td>Single processor</td>
</tr>
<tr>
<td>SuperMIC</td>
<td>workq</td>
<td>3 days</td>
<td>20</td>
<td>34</td>
<td>128</td>
<td>Unpreemptable</td>
</tr>
<tr>
<td></td>
<td>checkpoint</td>
<td></td>
<td>20</td>
<td></td>
<td>360</td>
<td>Preemptable</td>
</tr>
</tbody>
</table>
Queue Characteristics

➢ “qstat -q” will give you more info on the queues

[fchen14@mike2 ~]$ qstat -q

server: mike3

<table>
<thead>
<tr>
<th>Queue</th>
<th>Memory</th>
<th>CPU</th>
<th>Time</th>
<th>Walltime</th>
<th>Node</th>
<th>Run</th>
<th>Que</th>
<th>Lm</th>
<th>State</th>
</tr>
</thead>
<tbody>
<tr>
<td>workq</td>
<td>--</td>
<td>--</td>
<td>72:00:00</td>
<td>128</td>
<td>31</td>
<td>6</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>mwfa</td>
<td>--</td>
<td>--</td>
<td>72:00:00</td>
<td>8</td>
<td>3</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>bigmem</td>
<td>--</td>
<td>--</td>
<td>48:00:00</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>lasigma</td>
<td>--</td>
<td>--</td>
<td>72:00:00</td>
<td>28</td>
<td>28</td>
<td>7</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>bigmemtb</td>
<td>--</td>
<td>--</td>
<td>48:00:00</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>priority</td>
<td>--</td>
<td>--</td>
<td>168:00:00</td>
<td>128</td>
<td>0</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>single</td>
<td>--</td>
<td>--</td>
<td>72:00:00</td>
<td>1</td>
<td>62</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>gpu</td>
<td>--</td>
<td>--</td>
<td>24:00:00</td>
<td>16</td>
<td>1</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>preempt</td>
<td>--</td>
<td>--</td>
<td>72:00:00</td>
<td>--</td>
<td>0</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>checkpt</td>
<td>--</td>
<td>--</td>
<td>72:00:00</td>
<td>128</td>
<td>31</td>
<td>137</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>admin</td>
<td>--</td>
<td>--</td>
<td>24:00:00</td>
<td>--</td>
<td>0</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
<tr>
<td>scalemp</td>
<td>--</td>
<td>--</td>
<td>24:00:00</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>--</td>
<td>E</td>
<td>R</td>
</tr>
</tbody>
</table>

➢ For a more detailed description use mdiag
Queue Querying – Linux Clusters

- Displays information about active, eligible, blocked, and/or recently completed jobs: `showq` command

$ showq

active jobs------------------------
<table>
<thead>
<tr>
<th>JOBID</th>
<th>USERNAME</th>
<th>STATE</th>
<th>PROCs</th>
<th>REMAINING</th>
<th>STARTTIME</th>
</tr>
</thead>
<tbody>
<tr>
<td>236875</td>
<td>ebeigi3</td>
<td>Running</td>
<td>16</td>
<td>1:44:29</td>
<td>Mon Sep 15 20:00:22</td>
</tr>
<tr>
<td>236934</td>
<td>mwu3</td>
<td>Running</td>
<td>16</td>
<td>00:03:27</td>
<td>Mon Sep 15 19:04:20</td>
</tr>
</tbody>
</table>

eligible jobs----------------------
<table>
<thead>
<tr>
<th>JOBID</th>
<th>USERNAME</th>
<th>STATE</th>
<th>PROCs</th>
<th>WCLIMIT</th>
<th>QUEUETIME</th>
</tr>
</thead>
<tbody>
<tr>
<td>236795</td>
<td>dmarce1</td>
<td>Idle</td>
<td>1456</td>
<td>00:15:00</td>
<td>Mon Sep 15 16:38:45</td>
</tr>
<tr>
<td>236753</td>
<td>rsmith</td>
<td>Idle</td>
<td>2000</td>
<td>4:00:00</td>
<td>Mon Sep 15 14:44:52</td>
</tr>
<tr>
<td>236862</td>
<td>dlamas1</td>
<td>Idle</td>
<td>576</td>
<td>2:00:00</td>
<td>Mon Sep 15 17:28:57</td>
</tr>
</tbody>
</table>

121 eligible jobs

blocked jobs-----------------------
<table>
<thead>
<tr>
<th>JOBID</th>
<th>USERNAME</th>
<th>STATE</th>
<th>PROCs</th>
<th>WCLIMIT</th>
<th>QUEUETIME</th>
</tr>
</thead>
<tbody>
<tr>
<td>232741</td>
<td>myagho1</td>
<td>Idle</td>
<td>2000</td>
<td>1:00:00:00</td>
<td>Mon Sep 8 07:22:12</td>
</tr>
<tr>
<td>235545</td>
<td>tanping</td>
<td>Idle</td>
<td>1</td>
<td>2:21:10:00</td>
<td>Fri Sep 12 16:50:49</td>
</tr>
<tr>
<td>235546</td>
<td>tanping</td>
<td>Idle</td>
<td>1</td>
<td>2:21:10:00</td>
<td>Fri Sep 12 16:50:50</td>
</tr>
</tbody>
</table>

...
HPC User Environment 2

Submit and Monitor Your Jobs
Two Job Types

➢ Interactive job
  – Set up an interactive environment on compute nodes for users
    • Advantage: can run programs interactively
    • Disadvantage: must be present when the job starts
  – Purpose: testing and debugging, compiling
    • Do not run on the head node!!!
    • Try not to run interactive jobs with large core count, which is a waste of resources)

➢ Batch job
  – Executed without user intervention using a job script
    • Advantage: the system takes care of everything
    • Disadvantage: can only execute one sequence of commands which cannot changed after submission
  – Purpose: production run
Submitting Jobs on Linux Clusters

➢ Interactive job example:

```bash
qsub -I -X \ 
-l walltime=<hh:mm:ss>,nodes=<num_nodes>:ppn=<num_cores> \ 
-A <Allocation> \ 
-q <queue name>
```

DO NOT directly `ssh` to compute nodes, unless the nodes are assigned to you by the job scheduler.

– Add `-X` to enable X11 forwarding

➢ Batch Job example:

```bash
qsub job_script
```
Check Your Available Allocations

[fchen14@mike2 ~]$ balance

=================================== Allocation information for fchen14 ======================

<table>
<thead>
<tr>
<th>Proj. Name</th>
<th>Alloc</th>
<th>Balance</th>
<th>Deposited</th>
<th>%Used</th>
<th>Days Left</th>
<th>End</th>
</tr>
</thead>
<tbody>
<tr>
<td>hpc_hpcadmin3</td>
<td>hpc_hpcadmin3 on @mike2</td>
<td>282854.91</td>
<td>350000.00</td>
<td>19.18</td>
<td></td>
<td>2017-06-30</td>
</tr>
<tr>
<td>hpc_trn17mike2</td>
<td>hpc_trn17mike2 on @mike2</td>
<td>20305.62</td>
<td>25000.00</td>
<td>18.78</td>
<td></td>
<td>2018-04-01</td>
</tr>
</tbody>
</table>

Note: Balance and Deposit are measured in CPU-hours

[fchen14@mike2 ~]$ showquota

Hard disk quotas for user fchen14 (uid 32584):

<table>
<thead>
<tr>
<th>Filesystem</th>
<th>MB used</th>
<th>quota</th>
<th>files</th>
<th>fquota</th>
</tr>
</thead>
<tbody>
<tr>
<td>/home</td>
<td>4518</td>
<td>5000</td>
<td>94354</td>
<td></td>
</tr>
<tr>
<td>/work</td>
<td>424228</td>
<td>0</td>
<td>286002</td>
<td>400000</td>
</tr>
<tr>
<td>/project</td>
<td>65346</td>
<td>100000</td>
<td>1119432</td>
<td>400000</td>
</tr>
</tbody>
</table>

CPU Allocation SUs remaining:

- hpc_hpcadmin3: 282854.91
- hpc_trn17mike2: 20305.62
Submit An Interactive Job on SuperMike2

[fchen14@mike1 ~]$ qsub -I -X -l nodes=1:ppn=16,walltime=2:00:00 -q workq -A hpc_train_2018

qsub: waiting for job 675733.mike3 to start
qsub: job 675733.mike3 ready

Running PBS prologue script
...
Job ID: 675733.mike3
Username: fchen14
Group: Admins
Date: 13-Jun-2018 15:34
Node: mike044 (62703)

PBS has allocated the following nodes:
mike044
A total of 16 processors on 1 nodes allocated

Concluding PBS prologue script - 13-Jun-2017 15:34:19

[fchen14@mike044 ~]$
PBS Environmental Variables

[fchen14@mike315 ~]$ echo $PBS_

<table>
<thead>
<tr>
<th>Environment Variable</th>
</tr>
</thead>
<tbody>
<tr>
<td>$PBS_ENVIRONMENT</td>
</tr>
<tr>
<td>$PBS_MOMPORT</td>
</tr>
<tr>
<td>$PBS_NUM_PPN</td>
</tr>
<tr>
<td>$PBS_O_MAIL</td>
</tr>
<tr>
<td>$PBS_QUEUE</td>
</tr>
<tr>
<td>$PBS_WALLTIME</td>
</tr>
<tr>
<td>$PBS_GPUFILE</td>
</tr>
<tr>
<td>$PBS_NODEFILE</td>
</tr>
<tr>
<td>$PBS_O_HOME</td>
</tr>
<tr>
<td>$PBS_O_PATH</td>
</tr>
<tr>
<td>$PBS_SERVER</td>
</tr>
<tr>
<td>$PBS_JOBCOOKIE</td>
</tr>
<tr>
<td>$PBS_NODENUM</td>
</tr>
<tr>
<td>$PBS_O_HOST</td>
</tr>
<tr>
<td>$PBS_O_QUEUE</td>
</tr>
<tr>
<td>$PBS_TASKNUM</td>
</tr>
<tr>
<td>$PBS_JOBID</td>
</tr>
<tr>
<td>$PBS_NP</td>
</tr>
<tr>
<td>$PBS_O_LANG</td>
</tr>
<tr>
<td>$PBS_O_SHELL</td>
</tr>
<tr>
<td>$PBS_VERSION</td>
</tr>
<tr>
<td>$PBS_JOBNAME</td>
</tr>
<tr>
<td>$PBS_NUM_NODES</td>
</tr>
<tr>
<td>$PBS_O_LOGNAME</td>
</tr>
<tr>
<td>$PBS_O_WORKDIR</td>
</tr>
<tr>
<td>$PBS_VNODENUM</td>
</tr>
</tbody>
</table>
#PBS Job Script – Serial Job

#!/bin/bash
#PBS -l nodes=1:ppn=1      # Number of nodes and processor
#PBS -l walltime=24:00:00  # Maximum wall time
#PBS -N myjob             # Job name
#PBS -o <file name>       # File name for standard output
#PBS -e <file name>       # File name for standard error
#PBS -q single            # The queue for serial jobs
#PBS -A <loni_allocation> # Allocation name
#PBS -m e                 # Send mail when job ends
#PBS -M <email address>   # Send mail to this address

<shell commands>
<path_to_executable> <options>
<shell commands>

Tells the job scheduler how much resource you need.

How will you use the resources?
PBS Job Script – Parallel Job

#!/bin/bash

#PBS -l nodes=2:ppn=16  # Number of nodes and processors per node
#PBS -l walltime=24:00:00 # Maximum wall time
#PBS -N myjob            # Job name
#PBS -o <file name>      # File name for standard output
#PBS -e <file name>      # File name for standard error
#PBS -q checkpt          # Queue name
#PBS -A <allocation_if_needed> # Allocation name
#PBS -m e               # Send mail when job ends
#PBS -M <email address>  # Send mail to this address

<shell commands>

mpirun -machinefile $PBS_NODEFILE -np 32 <path_to_executable> <options>

<shell commands>
True or False?

➢ I have the below job script on QB2, since I used nodes=2:ppn=20, my script will run in parallel using 2 nodes with 40 cores.

   a) True
   b) False

#!/bin/bash
#PBS -l nodes=2:ppn=20
#PBS -l walltime=24:00:00
#PBS -N myjob
#PBS -j oe
#PBS -q checkpt
#PBS -A my_allocation

./my_executable.out
Job Monitoring - Linux Clusters

- Check details on your job using `qstat`
  
  ```
  $ qstat -n -u $USER : For quick look at nodes assigned to you
  $ qstat -f jobid : For details on your job
  $ qdel jobid : To delete job
  ```

- Check approximate start time using `showstart`
  
  ```
  $ showstart jobid
  ```

- Check details of your job using `checkjob`
  
  ```
  $ checkjob jobid
  ```

- Check health of your job using `qshow`
  
  ```
  $ qshow jobid
  ```

- Please pay close attention to the load and the memory consumed by your job!
Using the “top” command

➢ The top program provides a dynamic real-time view of a running system.

```
top - 19:39:56 up 89 days, 4:13, 1 user, load average: 0.63, 0.18, 0.06
Tasks: 489 total, 2 running, 487 sleeping, 0 stopped, 0 zombie
Cpu(s): 6.3%us, 0.0%sy, 0.0%ni, 93.7%id, 0.0%wa, 0.0%hi, 0.0%si, 0.0%st
Mem: 65909356k total, 3389616k used, 62519740k free, 151460k buffers
Swap: 207618040k total, 5608k used, 207612432k free, 947716k cached

PID USER      PR NI VIRT  RES  SHR S %CPU %MEM    TIME+ COMMAND
39595 fchen14   20   0 266m 257m  592 R 99.9  0.4   0:06.94 a.out
39589 fchen14   20   0 17376 1612  980 R  0.3  0.0   0:00.05 top
38479 fchen14   20   0 108m 2156 1348 S  0.0  0.0   0:00.03 bash
39253 fchen14   20   0 103m 1340 1076 S  0.0  0.0   0:00.00 236297.mike3.SC
39254 fchen14   20   0 103m 1324 1060 S  0.0  0.0   0:00.00 bm_laplace.sh
39264 fchen14   20   0 99836 1908  992 S  0.0  0.0   0:00.00 sshd
39265 fchen14   20   0 108m 3056 1496 S  0.0  0.0   0:00.03 bash
```
Pay attention to single queue usage

➢ Single queue - Used for jobs that will only execute on a single node, i.e. *nodes=1:ppn=1/2/4/6/8*.

➢ Jobs in the single queue should not use:
  – More than 2GB memory per core for Eric, Philip and SuperMike2 (32G/16).
  – More than 3.2GB memory per core for QB2 (64G/20).

➢ If applications require more memory, scale the number of cores (ppn) to the amount of memory required: i.e. max memory available for jobs in single queue is 8GB for pnn=4 on SuperMikell.

➢ Typical type of warning:
  – E124 - *Exceeded memory allocation*. This Job XXXX appears to be using more memory (GB) than allocated (9 > 3).
  – E123 - *Exceeded pnn/core allocation*. This Job XXXX appears to be using more cores than allocated (6 > 1). Please allocate the number of cores that the job will use, (ppn=6). This Job has 1 core(s) allocated (ppn=1).
Core and Memory in Single queue

20 cores

64GB memory

\[ \frac{64}{20} = 3.2 \text{GB} \]

**Question:**
On QB2, if my job needs 7GB memory, what ppn value should I use?
On SuperMike2, if my job needs 7GB memory, what ppn value should I use?
More things to be noticed

- Eric is old and will be retired in the near future. LONI users are encouraged to migrate their codes to QB-2 as soon as possible.
- The purpose of bigmem queue on QB-2 is for jobs costing big (larger than 64 GB) memory not for jobs using more number of cores.
- GPU is available to workq or checkpt queues on QB-2.
- Xeon Phi is available to workq or checkpt queues on SuperMIC.
- There is no single queue on SuperMIC.
- Users are encouraged to use accelerators (GPU/Xeon Phi) whenever possible. Application for allocation involving with usage of accelerators will be easier to be approved.
Job Submission Quiz

➢ How to suspend your account? (cont’d)
   – Use more memory than allowed. (e.g. use 5GB memory on SuperMike2 with ppn=1)
   – Seriously underutilize node resources (e.g. allocate 32 nodes but just use 1 core)
   – Submit job to the big memory queue but use only few MB of memory

➢ How to monitor core and memory usage?
Job Scheduling Basics
As a user, you interact with the scheduler and/or resource manager whenever you submit a job, or query on the status of your jobs or the whole cluster, or seek to manage your jobs.

Resource managers give access to compute resource
- Takes in a resource request (job) on login node
- Finds appropriate resource and assigns you a priority number
- Positions your job in a queue based on the priority assigned.
- Starts running jobs until it cannot run more jobs with what is available.
Job Scheduler

- HPC & LONI Linux clusters use TORQUE, an open source version of the Portable Batch System (PBS) together with the MOAB Scheduler, to manage user jobs.

- Resource Manager - Torque
  - Manages a queue of jobs for a cluster of resources
  - Launches job to a simple FIFO job queue

- Workload Manager - Moab
  - A scheduler that integrates with one or more Resource Managers to schedule jobs across domains of resources (servers, storage, applications)
  - Prioritizes jobs
  - Provides status of running and queued jobs, etc.

- The batch queuing system determines
  - The order jobs are executed
  - On which node(s) jobs are executed
Job management philosophy

➢ Working Philosophy
  – Prioritize workload into a queue for jobs
  – **Backfill** idle nodes to maximize utilization
    • Will be detailed later...
Job Priorities

- Jobs with a higher job priority are scheduled ahead of jobs with a lower priority.
- Job priorities have contributions from the following:
  - credential priority
  - fairshare priority
  - resource priority
  - service priority
- Priority determination for each queued job, use
  - `mdiag -p`:

```
$ mdiag -p
diagnosing job priority information (partition: ALL)
Job          PRIORITY*   Cred( User:Class)   FS( User: WCA)   Serv(QTime:XFctr)   Res( Proc)
Weights ------ 100( 10: 10) 100( 10: 50) 2( 2: 20) 30( 10)
---
236172        246376        40.6(100.0: 0.0) 8.6( 19.6: 0.3) 4.0(1480.: 99.7) 46.8(2048.)
235440        242365        41.3(100.0: 0.0) 4.6(  8.2: 0.6) 6.6(3959.:  6.5) 47.5(512.0)
235441        242365        41.3(100.0: 0.0) 4.6(  8.2: 0.6) 6.6(3959.:  6.5) 47.5(512.0)
235442        242361        41.3(100.0: 0.0) 4.6(  8.2: 0.6) 6.6(3958.:  6.5) 47.5(512.0)
236396        241821        41.4(100.0: 0.0) 8.8( 19.6: 0.3) 2.2(664.0: 67.4) 47.6(1456.)
```
Priority components

➢ **Credential priority**  
\[ \text{credweight} \times (\text{userweight} \times \text{job.user.priority}) \]
\[ = 100 \times (10 \times 100) = 100000 \]
It is a **constant** for all users.

➢ **Fairshare priority**  
\[ \text{fsweight} \times \text{min}(\text{fscap}, (\text{fsuserweight} \times \text{DeltaUserFSUsage})) \]
\[ = 100 \times (10 \times \text{DeltaUserFSUsage}) \]
If you have not submitted jobs in the past 7 days, \( \text{DeltaUserFSUsage} = 20000 \)

➢ **Service priority**  
\[ \text{serviceweight} \times (\text{queuetimeweight} \times \text{QUEUETIME} + \text{xfactorweight} \times \text{XFACTOR}) \]
\[ = 2 \times (2 \times \text{QUEUETIME} + 20 \times \text{XFACTOR}) \],
where \( \text{XFACTOR} = 1 + \frac{\text{QUEUETIME}}{\text{WALLTIMELIMIT}} \).

➢ **Resource priority**  
\[ \text{resweight} \times \text{min}(\text{rescap}, (\text{procweight} \times \text{TotalProcessorsRequested})) \]
\[ = 30 \times \text{min}(3840, (10 \times \text{TotalProcessorsRequested})) \]

➢ See [http://www.hpc.lsu.edu/docs/pbs.php](http://www.hpc.lsu.edu/docs/pbs.php), click “Job priority”. 

How to get higher priority?

➢ Do not submit too many jobs within one week.
➢ Submit your job early to accumulate the queue time.
➢ More on resource priority:
  – Request more compute nodes.
  – Request a smaller walltime limit.
  – see next few slides...
How to maximize the usage of a cluster?

- Fill in high-priority (large) jobs
- Backfill low-priority (small) jobs
An Overview of Backfilling (1)

➢ Backfill is a scheduling optimization that allows a scheduler to make better use of available resources by running jobs out of order.

➢ Enabling backfill allows the scheduler to start other, lower-priority jobs so long as they do not delay the highest priority job.

➢ If the FIRSTFIT algorithm is applied, the following steps are taken:
  – The list of feasible backfill jobs is filtered, selecting only those that will actually fit in the current backfill window.
  – The first job is started.
  – While backfill jobs and idle resources remain, repeat step 1.
An Overview of Backfilling (2)

➢ Although by default the start time of the highest priority job is protected by a reservation, there is nothing to prevent the third priority job from starting early and possibly delaying the start of the second priority job.

➢ Command to show current backfill windows:
  - showbf
    • Shows what resources are available for immediate use.
    • This command can be used by any user to find out how many processors are available for immediate use on the system. It is anticipated that users will use this information to submit jobs that meet these criteria and thus obtain quick job turnaround times.
  - Example:

```
[fchen14@eric2 ~]$ showbf -c workq
Partition   Tasks  Nodes   Duration     StartOffset     StartDate
---------   -----  -----   -----------     --------------     -----------
ALL        40      5 18:50:35     00:00:00      11:16:49_09/04
ALL         8      1 INFINITY     00:00:00      11:16:49_09/04
```
How Much Time Should I Ask for?

➢ **It should be**
  - Long enough for your job to complete
  - As short as possible to increase the chance of backfilling
Frequently Asked Questions

➢ I submitted job A before job B. Why job B started earlier than job A?
➢ There are free nodes available, why my job is still waiting and not running?
➢ Why my job is not get accelerated when running on cluster?
  – Is your job utilizing the parallel resource on the cluster?
  – Does you job have lots of I/O tasks?
  – See next section...
Compile and Analyze Codes on Cluster
Compilers

➢ Serial compilers

<table>
<thead>
<tr>
<th>Language</th>
<th>Linux cluster</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Intel</td>
</tr>
<tr>
<td>Fortran</td>
<td>ifort</td>
</tr>
<tr>
<td>C</td>
<td>icc</td>
</tr>
<tr>
<td>C++</td>
<td>icpc</td>
</tr>
</tbody>
</table>

➢ Parallel compilers

<table>
<thead>
<tr>
<th>Language</th>
<th>Linux clusters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fortran</td>
<td>mpif77, mpif90</td>
</tr>
<tr>
<td>C</td>
<td>mpicc</td>
</tr>
<tr>
<td>C++</td>
<td>mpiCC</td>
</tr>
</tbody>
</table>
Example compiling serial code

- `icc hello_cpu_elapsed.c`
- `gfortran test_hello2.f90`

- **List symbols for executables:**
  - `nm - list symbols from object files`
  - **Example:**
    ```
    [fchen14@mike2 hello]$ nm ./a.out | grep intel
    000000000060eb60 B __intel_cpu_indicator
    
    [fchen14@mike2 hello]$ nm ./a.out | grep gfortran
    U _gfortran_set_args@@GFORTRAN_1.0
    ```
CPU time vs Elapsed time

➢ **CPU time (or process time):**
  - The amount of time for which a central processing unit (CPU) was used for processing instructions of a computer program or operating system, as opposed to, for example, waiting for input/output (I/O) operations or entering low-power (idle) mode.

➢ **Elapsed real time (or simply real time, or wall clock time)**
  - The time taken from the start of a computer program until the end as measured by an ordinary clock. Elapsed real time includes I/O time and all other types of waits incurred by the program.

➢ **If a program uses parallel processing, total CPU time for that program would be more than its elapsed real time.**
  - \( \frac{\text{Total CPU time}}{\text{Number of CPUs}} \) would be same as elapsed real time if work load is evenly distributed on each CPU and no wait is involved for I/O or other resources.
#include <stdio.h>
#include <time.h>

int main(char *argc, char **argv) {
    double s=0.0;
    // fundamental arithmetic type representing clock tick counts.
    clock_t start, end;
    int i;
    start = clock();
    for (i=0;i<1000000000;i++)
        s+=i*2.0; // doing some floating point operations
    end = clock();
    double time_elapsed_in_seconds = (end - start)/(double)CLOCKS_PER_SEC;
    printf("cputime_in_sec: %e\n", time_elapsed_in_seconds);
    start = clock();
    system("sleep 5"); // just sleep, does this accumulate CPU time?
    end = clock();
    time_elapsed_in_seconds = (end - start)/(double)CLOCKS_PER_SEC;
    printf("cputime_in_sec: %e\n", time_elapsed_in_seconds);
    return 0;
}
Watch the actual cpu time using “time”

[fchen14@mike429 serial]$ gcc hello_cpu_elapsed.c
[fchen14@mike429 serial]$ time ./a.out

cputime_in_sec:  2.740000e+00
cputime_in_sec:  0.000000e+00

real  0m7.782s
user  0m2.750s
sys   0m0.005s
Some additional info about “time”

➢ Use the Linux command `time` to evaluate the actual time usage
  – time a simple command or give resource usage

➢ Real refers to actual elapsed time (wall clock time)
  – Time from start to finish of the call. This is all elapsed time including time used by other processes and time the process spends blocked (for example if it is waiting for I/O to complete).

➢ User and Sys refer to CPU time used only by the process.
  – User is the amount of CPU time spent in user-mode code (`outside the kernel`) within the process.
  – Sys is the amount of CPU time spent `in the kernel` within the process.

➢ Purpose of this example:
  – real < user: The process is CPU bound and takes advantage of parallel execution on multiple cores/CPUs.
  – real ≈ user: The process is CPU bound and takes no advantage of parallel execution.
  – real > user: The process is I/O bound. Execution on multiple cores would be of little to no advantage.
Two parallel schemes

➢ **Shared Memory system**
  – A single multicore compute node
  – Open Multi-processing (OpenMP)

➢ **Distributed Memory system**
  – Multiple compute nodes
  – Message Passing Interface (MPI)

**MPI**: Distributed Memory System

**OpenMP**: Shared Memory System

Typically less memory overhead/duplication. Communication often implicit, through cache coherency and runtime.
Example compiling threaded OpenMP code

- **Compiling OpenMP code often requires the openmp compiler flags, it varies with different compiler**
- **Environment Variable** `OMP_NUM_THREADS` **sets the number of threads**
- **Examples:**
  
  ```
  [fchen14@mike2 src]$ gcc -fopenmp hello_openmp.c
  [fchen14@mike2 src]$ ifort -openmp hello_openmp.f90
  ```

<table>
<thead>
<tr>
<th>Compiler</th>
<th>Compiler Options</th>
<th>Default behavior for # of threads (OMP_NUM_THREADS not set)</th>
</tr>
</thead>
<tbody>
<tr>
<td>GNU (gcc, g++, gfortran)</td>
<td>-fopenmp</td>
<td>as many threads as available cores</td>
</tr>
<tr>
<td>Intel (icc ifort)</td>
<td>-openmp</td>
<td>as many threads as available cores</td>
</tr>
<tr>
<td>Portland Group</td>
<td>-mp</td>
<td>one thread</td>
</tr>
<tr>
<td>(pgcc,pgCC,pgf77,pgf90)</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
#include <omp.h>
#include <stdio.h>
#include <stdlib.h>

int main (int argc, char *argv[]) {
    int nthreads, tid;
    /* Fork a team of threads with their own copies of variables */
    #pragma omp parallel private(nthreads, tid)
    {
        /* Obtain thread number */
        tid = omp_get_thread_num();
        printf("Hello World from thread = %d\n", tid);
        /* Only master thread does this */
        if (tid == 0) {
            nthreads = omp_get_num_threads();
            printf("Number of threads = %d\n", nthreads);
        }
    } /* All threads join master thread and disband */
}
Sample OpenMP Fortran code

```fortran
program hello

integer nthreads,tid,omp_get_num_threads,omp_get_thread_num
! fork a team of threads giving them their own copies of variables
!$omp parallel private(nthreads, tid)
! obtain thread number
  tid = omp_get_thread_num()
print *, 'hello world from thread = ', tid
! only master thread does this
if (tid .eq. 0) then
  nthreads = omp_get_num_threads()
  print *, 'number of threads = ', nthreads
end if
! all threads join master thread and disband
!$omp end parallel
end
```
Analyzing a parallel (OpenMP) program

➢ What will be the CPU time and elapsed time for the following code segment:

See (on SuperMike II):
/home/fchen14/userenv/src/openmp/hello_openmp_cpu_elapse.c

```c
// fundamental arithmetic type representing clock tick counts.
clock_t start, end;
struct timeval r_start, r_end;
int i;
gettimeofday(&r_start, NULL);
start = clock();
#pragma omp parallel for // spawn the openmp threads
for (i=0;i<N;i++) a = i*2.0; // doing some floating point operations
end = clock();
gettimeofday(&r_end, NULL);
double cputime_elapsed_in_seconds = (end - start)/(double)CLOCKS_PER_SEC;
double realtime_elapsed_in_seconds = ((r_end.tv_sec * 1000000 +
r_end.tv_usec) - (r_start.tv_sec * 1000000 +
r_start.tv_usec))/1000000.0;
```
# Available MPI libraries on LONI & HPC

<table>
<thead>
<tr>
<th>Cluster Resource</th>
<th>Name</th>
<th>MPI Library</th>
<th>Default serial</th>
<th>compiler</th>
</tr>
</thead>
<tbody>
<tr>
<td>LONI</td>
<td>Eric</td>
<td>Mvapich 0.98, 1.1</td>
<td>1.3.4</td>
<td>X</td>
</tr>
<tr>
<td></td>
<td>QB2</td>
<td>X</td>
<td>1.8.1</td>
<td>3.0.3</td>
</tr>
<tr>
<td>LSU</td>
<td>SuperMikell</td>
<td>X 1.9, 2.0.1</td>
<td>1.6.2</td>
<td>3.0.2</td>
</tr>
<tr>
<td></td>
<td>Philip</td>
<td>X</td>
<td>1.4.3, 1.6.1</td>
<td>1.2.7, 1.3.2, 1.4.1</td>
</tr>
<tr>
<td></td>
<td>SuperMIC</td>
<td>X 2.0</td>
<td>1.8.1</td>
<td>3.0.3, 3.1.1</td>
</tr>
</tbody>
</table>
MPI Compilers (1)

<table>
<thead>
<tr>
<th>Language</th>
<th>Linux clusters</th>
<th>AIX clusters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fortran</td>
<td>mpif77, mpif90</td>
<td>mpxlf, mpxlf90</td>
</tr>
<tr>
<td>C</td>
<td>mpicc</td>
<td>mpcc</td>
</tr>
<tr>
<td>C++</td>
<td>mpiCC</td>
<td>mpCC</td>
</tr>
</tbody>
</table>

mpif90 hello.f90

mpicc hello.c

mpicxx hello.cpp
MPI Compilers (2)

➢ These MPI compilers are actually wrappers
  – They still use the compilers we've seen on the previous slide
    • Intel, PGI or GNU
  – They take care of everything we need to build MPI codes
    • Head files, libraries etc.
  – What they actually do can be revealed by the -show option

➢ It’s extremely important that you compile and run your code with the same version of MPI!
  – Use the default version if possible
Compiling a MPI C program

- Compiling Hello world in C version:
  ```c
  mpicc hello_mpi.c
  #include <mpi.h>
  #include <stdio.h>
  int main(int argc, char** argv) {
      int name_len, world_size, world_rank;
      char processor_name[MPI_MAX_PROCESSOR_NAME];
      // Initialize the MPI environment
      MPI_Init(NULL, NULL);
      // Get the number and rank of processes
      MPI_Comm_size(MPI_COMM_WORLD, &world_size);
      MPI_Comm_rank(MPI_COMM_WORLD, &world_rank);
      // Get the name of the processor
      MPI_Get_processor_name(processor_name, &name_len);
      // Print off a hello world message
      printf("Iam from processor %s, rank %d out of %d processors\n", 
             processor_name, world_rank, world_size);
      // Finalize the MPI environment.
      MPI_Finalize();
  }
  ```
Compiling a MPI Fortran program

➢ Compiling Hello world in Fortran:
  – mpif90 hellp_mpi.f90

```fortran
program hello_mpi

  include 'mpif.h'
  !use mpi
  character 10 name
  ! Initialize the MPI library:
  call MPI_Init(ierr)
  ! Get size and rank
  call MPI_Comm_Size(MPI_COMM_WORLD, numtasks, ierr)
  call MPI_Comm_Rank(MPI_COMM_WORLD, rank, ierr)
  ! print date
  if (nrank == 0) then
      write( , )'System date'
      call system('date')
  endif
  call MPI_Barrier(MPI_COMM_WORLD, ierr)
  ! print rank
  call MPI_Get_Processor_Name(name, len, ierr)
  write( , )"I am ", nrank, "of", numtasks, "on ", name
  ! Tell the MPI library to release all resources it is using:
  call MPI_Finalize(ierr)
end program hello_mpi
```

02/21/2018

HPC User Environment 2 Fall 2018
Notes for compiling a MPI program (1)

➢ Always verify what compiler/library is being used:

$ mpicc -show
icc -I/usr/local/packages/openmpi/1.6.2/Intel-13.0.0/include -L/usr/local/packages/openmpi/1.6.2/Intel-13.0.0/lib -lmpi -ldl -lm -Wl,--export-dynamic -lrt -lnc -libverbs -libumad -lpthread -lutil

$ mpif90 -show
ifort -I/usr/local/packages/openmpi/1.6.2/Intel-13.0.0/include -I/usr/local/packages/openmpi/1.6.2/Intel-13.0.0/lib -L/usr/local/packages/openmpi/1.6.2/Intel-13.0.0/lib -lmpi_f90 -lmpi_f77 -lmpi -ldl -lm -Wl,--export-dynamic -lrt -lnc -libverbs -libumad -lpthread -lutil
Notes for compiling a MPI program (2)

➢ Always verify what library is being used: Before and after:

$ ldd a.out # ldd - print shared library dependencies

    linux-vdso.so.1 => (0x00007fff907ff000)
    libmpi_f90.so.1 => /usr/local/packages/openmpi/1.6.2/Intel-13.0.0/lib/libmpi_f90.so.1 (0x00002b9ae577e000)
    libmpi_f77.so.1 => /usr/local/packages/openmpi/1.6.2/Intel-13.0.0/lib/libmpi_f77.so.1 (0x00002b9ae5982000)
    libmpi.so.1 => /usr/local/packages/openmpi/1.6.2/Intel-13.0.0/lib/libmpi.so.1 (0x00002b9ae5bb9000)

    ... 
    libpthread.so.0 => /lib64/libpthread.so.0 (0x00000003b2180000)
    ...

    libifport.so.5 => /usr/local/compilers/Intel/composer_xe_2013.0.079/compiler/lib/intel64/libifport.so.5 (0x00002b9ae61ee000)
    libifcore.so.5 => /usr/local/compilers/Intel/composer_xe_2013.0.079/compiler/lib/intel64/libifcore.so.5 (0x00002b9ae641d000)
Running and Analyzing MPI program

➢ Make sure you are running your jobs on the correct nodes
➢ Important if you want to run less processes than ppn
➢ Understand the usage of $PBS_NODEFILE

[fchen14@mike2 ~]$ qsub -I -X -l nodes=2:ppn=16 -l walltime=01:00:00 -q gpu
...
[fchen14@mike429 ~]$ echo $PBS_NODEFILE
/var/spool/torque/aux/236660.mike3
[fchen14@mike429 ~]$ cat $PBS_NODEFILE
mike429
...  # 16 repeats of mike429
mike429
mike430
...  # 16 repeats of mike430
mike430
[fchen14@mike429 hybrid]$ cat $PBS_NODEFILE| uniq > hosts
[fchen14@mike429 hybrid]$ cat hosts
mike429
mike430
Running and Analyzing MPI program

[fchen14@mike315 mpi]$ mpicc hello_mpi.c
[fchen14@mike315 mpi]$ mpirun -np 32 -hostfile $PBSNODEFILE ./a.out

Iam from processor mike315, rank 1 out of 32 processors
Iam from processor mike315, rank 6 out of 32 processors
Iam from processor mike315, rank 9 out of 32 processors
Iam from processor mike315, rank 12 out of 32 processors
Iam from processor mike315, rank 0 out of 32 processors
Iam from processor mike315, rank 2 out of 32 processors
Iam from processor mike315, rank 3 out of 32 processors
Iam from processor mike315, rank 7 out of 32 processors
Iam from processor mike315, rank 10 out of 32 processors
Iam from processor mike315, rank 5 out of 32 processors
Iam from processor mike315, rank 13 out of 32 processors
Iam from processor mike315, rank 4 out of 32 processors
Iam from processor mike315, rank 8 out of 32 processors
Iam from processor mike334, rank 17 out of 32 processors
Iam from processor mike315, rank 11 out of 32 processors
Iam from processor mike315, rank 14 out of 32 processors
Iam from processor mike315, rank 15 out of 32 processors
Iam from processor mike334, rank 18 out of 32 processors
Compiling hybrid (MPI+OpenMP) program

➢ See /home/fchen14/userenv/src/hybrid/hello_hybrid.c for complete source
➢ Use command:
  - $ mpicc -openmp hello_hybrid.c

#pragma omp parallel default(shared) private(itd, np)
{
    gtd = omp_get_num_threads(); // get total num of threads in a process
    itd = omp_get_thread_num();  // get thread id
    gid = nrank*gtd + itd;       // global id
    printf("Gid %d from thd %d out of %d from process %d out of %d on %s\n", 
        gid, itd, gtd, nrank, numprocs, processor_name);
    if (nrank==0 && itd==0)
    {
        // system("pstree -ap -u $USER");
        system("for f in `cat $PBS_NODEFILE|uniq`; do ssh $f pstree -ap -u $USER; done;"));
        system("sleep 10");
    }
}
Analyzing a hybrid program

[fchen14@mike315 hybrid]$ export OMP_NUM_THREADS=4
[fchen14@mike315 hybrid]$ mpirun -np 2 -x OMP_NUM_THREADS ./a.out
Gid 0 from thread 0 out of 4 from process 0 out of 2 on mike315
Gid 2 from thread 2 out of 4 from process 0 out of 2 on mike315
Gid 1 from thread 1 out of 4 from process 0 out of 2 on mike315
Gid 3 from thread 3 out of 4 from process 0 out of 2 on mike315
Gid 4 from thread 0 out of 4 from process 1 out of 2 on mike315
Gid 6 from thread 2 out of 4 from process 1 out of 2 on mike315
Gid 7 from thread 3 out of 4 from process 1 out of 2 on mike315
Gid 5 from thread 1 out of 4 from process 1 out of 2 on mike315
bash,108067
  | -mpirun,110651 -np 2 -x OMP_NUM_THREADS ./a.out
  |   | -a.out,110652
  |   |   | -sh,110666 -c ...
  |   |   | `ssh,110670 mike315 pstree -ap -u fchen14
  |   |   | `-a.out,110654
  |   |   | `-a.out,110656
  |   |   | `-a.out,110662
  |   |   | `-a.out,110663
  |   |   | `-a.out,110664
  |   |   | `-{a.out},110665
  |
Exercise

➢ Submit a small job to run “sleep 180” and “print PBS variables”
  – Create a script to submit a 5 min job and print from within the job script
    PBS variables $PBS_NODEFILE, $PBS_WORKDIR. Also run “sleep 180” to give you a few minutes to verify status.
  – Once the job is running, find out the Mother Superior node and other
    slave nodes assigned to your job using qstat.
  – Log into MS node and verify that your job is running and find your
    temporary output file
  – Modify your script to print hello from each of your assigned nodes

➢ Run a shell script using mpirun to print process id of shell
Future Trainings

➢ Next week training: Basic Shell Scripting
   – Wednesdays 9:00am, September 20, Frey Computing Service Center 307

➢ Programming/Parallel Programming workshops
   – Usually in summer

➢ Keep an eye on our webpage: www.hpc.lsu.edu