

Parallelizing The Matrix Multiplication



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Serial version



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$$c_{i,j} = \sum_{k=1}^d a_{i,k} \cdot b_{k,j}$$







//Initialize the arrays
For all elements of A and B
initial value = function (i , j)

// Matrix multiplication
For each C (i,j)
// Take the inner product of row i of A and column j of B
For each A (i,k) and B(k,j)
C (i,j) = C (i,j) + A (i,k) * B(k,j)

 $a_{i,j} = i + j$ $b_{i,j} = i * j$





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 $\begin{array}{l}a_{i,j}=i+j\\b_{i,j}=i*j\end{array}$

Anything else?



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How do we know the result is correct? **Result validation is very important!**











$$c_{i,j} = \sum_{k=1}^{N} a_{i,k} \cdot b_{k,j}$$

$$= \sum_{k=1}^{N} (i+k) \cdot j \cdot k$$
Assuming A, B and C are all
N x N square matrices
$$= \sum_{k=1}^{N} (ijk+jk^2)$$

$$= ij \sum_{k=1}^{N} k+j \sum_{k=1}^{N} k^2$$

$$= ij \cdot \frac{N(N+1)}{2} + j \cdot \frac{N(N+1)(2N+1)}{6}$$

We can use this formula to validate the result from the program



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//Validate the result Print out an element of C and compare it to the value given by the formula







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Serial Program Details

- A, B and C are all N x N square matrices
- Take input arguments
 - N: dimension of the matrices
 - $-i_{peek}$, j_{peek} : indices of the element used for result validation
- Values of A and B:

$$a_{i,j} = i + j$$
$$b_{i,j} = i * j$$

- Output
 - How much time it takes (performance measurement)
 - The value of $C(i_{peek}, j_{peek})$ and the value given by the formula







Parallel version



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Writing a parallel program step by step

- Step 1. Start from serial programs as a baseline
 Something to check correctness and efficiency against
- Step 2. Analyze and profile the serial program
 - Identify the "hotspot"
 - Identify the parts that can be parallelized
- Step 3. Parallelize code incrementally
- Step 4. Check correctness of the parallel code
- Step 5. Iterate step 3 and 4







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Which parts can and should be parallelized?





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 $c_{i,j} = \sum_{k=1}^{n} a_{i,k} \cdot b_{k,j}$







Some Considerations

- Decomposition
 - 1-D or 2-D
- Data distribution
 - Each process owns the entire matrices, but only performs calculation on a part of them, or
 - Each process only owns the sub-matrices that it is going to process









 $c_{i,j} = \sum_{k=1}^{N} a_{i,k} \cdot b_{k,j}$













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Optimization Conderations

- Avoid data movement as much as possible (i.e. increase the amount of computation done relative to the amount of data moved)
 - True for both serial and parallel programs
 - For parallel programs, data movement means data communication between host and device (GPU) or between different nodes (distributed memory systems)
- Reduce the memory footprint
- When developing a parallel program, it's important to know what to expect in terms of speedup and scaling







Cannon's Algorithm





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Cannon's Algorithm

- Assume
 - the number of processes p is a perfect square
 - the matrices are N x N square
- Arrange the processes into a 2-D $\sqrt{p} \ge \sqrt{p}$ process grid
 - For each matrix, each process is assigned with a block of $N/\sqrt{p} \ge N/\sqrt{p}$











Working by elements

 $c_{i,j} = \sum_{k=1}^{N} a_{i,k} \cdot b_{k,j}$





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 $C(1,2) = A(1,0) \times B(0,2) + A(1,1) \times B(1,2) + A(1,2) \times B(2,2)$



Working by blocks



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 $C(1,2) = \underline{A(1,0) \times B(0,2)} + \underline{A(1,1) \times B(1,2)} + \underline{A(1,2) \times B(2,2)}$

Three phases: phase1 phase2 phase3 The multiplication is completed in \sqrt{p} phases CENTER FOR COMPUTATION



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Cannon's Algorithm

- Two stages
 - Skew the matrices so everything aligns properly
 - Shift row i of A by i columns to the left
 - Shift column j of B by j rows to the up
 - Shift and multiply
 - Each process calculate the local product and add into the accumulated sum
 - Shift A by 1 column to the left
 - Shift B by 1 row to the up
 - Repeat \sqrt{p} times
- All shifts wrap around (circular)









 Process P(1,2) owns
 After initialization:
 A(1,2), B(1,2), C(1,2)

 After skewing:
 A(1,0), B(0,2), C(1,2)

 Shifting once:
 A(1,1), B(1,2), C(1,2)

 Shifting twice:
 A(1,2), B(2,2), C(1,2)



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Cannon's Algorithm – Pseudo Code

For i = 0 to sqrt(p) - 1Shift A(i,:) to the left by i For j = 0 to sqrt(p) - 1Shift B(:,j) to the up by j For k = 0 to sqrt(p) - 1For i = 0 to sqrt(p) - 1 and j = 0 to sqrt(p) - 1 $C(i,j) += A(i,j) \times B(i,j)$ Shift A(i,:) to the left by 1 Shift B(:,j) to the up by 1



